# The BOL IR

# Draft of October 21, 2004

# **1** Introduction

BOL is a normalized extended  $\lambda$ -calculus that serves as the intermediate representation (IR) of the MOBY compiler. It has a weak, but simple, type system that serves as a guide for optimization and code generation. This report describes the dynamic and static semantics of BOL. It is meant to serve as documentation for MOBY compiler.

The following table summarizes the SML types used to represent BOL types and terms and where they are defined and described:

Туре	Module	Description	Section
var	BOL	BOL variables	5
exp	BOL	BOL expressions labeled by program points	4.1
term	BOL	unlabeled expressions	4.1
lambda	BOL	named function definition	4.2
rhs	BOL	right-hand-side of binding	4.3
primop	PrimOps	primitive operations	6
ppt	ProgPt	program points	
kind	BOLTypes	kind of a BOL type	2.1
ty	BOLTypes	BOL type	2.2 - 2.5
field	BOLTypes	field descriptor for BOL struct	2
c_prototype	BOLTypes	type of C function	2

# 2 BOL types

#### 2.1 Kinds

The BOL types are organized into a hierarchy by kind; there are four distinct kinds of BOL types:

- 1. Word kind (W) types are those that can be stored in a general-purpose machine register on the host processor.
- 2. Variable kind types (V) are those that can be assigned to a BOL variable.
- 3. Memory kind types (M) are those types that describe the layout of memory.
- 4. Type kind types (Type) include all types.

We use Kind = {W, V, M, Type} for the set of kinds and  $\kappa \in$  Kind. The kinds are ordered under set inclusion as follows:

$$\mathbf{W} \subset \mathbf{V} \subset \mathbf{M} \subset \mathbf{Type}$$

*i.e.*, a type of kind **W** also has kinds **V**, **M**, and **Type**. A *kind environment* KE : (Base  $\cup$  TyVar)  $\xrightarrow{\text{fin}}$  Kind maps base types and type variables to kinds. The mapping of base types is architecture (and compiler) specific. For example, the type of 64-bit integers (long) has kind **W** on 64-bit machines, but kind **V** on 32-bit machines.

#### 2.2 Kind W types

The following types have W kind and may be mapped to a general-purpose register:

#### T\_Any

a word-sized value of unknown type; we use the syntax any to denote this type.

T\_Bool

a boolean; we use the syntax **bool** to denote this type.

T\_Enum of {lo : word, hi : word}

a small integer (16-bit) in the range [lo, hi]. When lo is equal to hi, then the type is a *singleton* type. We write (lo..hi) to denote the type T\_Enum{lo, hi}.

T\_Integer

arbitrary precision integers (represented by a pointer); we use the syntax **integer** to denote this type.

#### T\_Wrap **of** ty

a wrapped value (the type argument will be one of: int, long, float, double, or extended). We use the syntax wrap( $\tau$ ) to denote the type T\_Wrap( $\tau$ ).

T\_Addr **of** ty

the address of memory with the given type. The memory is guaranteed to be outside the MOBY heap. We use the syntax  $\&\tau$  to denote the type T\_Addr( $\tau$ ).

T\_Ptr **of** ty

a pointer to memory with the given type. The memory may be in the MOBY heap. We use the syntax  $*\tau$  to denote the type T\_Ptr( $\tau$ ).

T\_PtrOrEnum **of** {ptrTy : ty, enumTy : ty}

a value that is either a pointer to memory with ptrTy type or is an enumeration (enumTy specifies the range). We use the syntax \*ptrTy + (lo..hi) to denote the type

```
T_PtrOrEnum{ptrTy, enumTy=T_Enum{lo, hi}}
```

- T\_FunPtr of {dom : ty list, rng : ty list}
   a function with the given domain and range.
- T\_CodePtr **of** {dom : ty list, rng : ty list}

the address of machine code for a function with the given domain and range. Values of this type are introduced as part of closure conversion.

```
T_ContPtr of ty list
```

a BOL continuation with the given argument types.

```
T_Label of ty list
```

the address of an internal fragment in a cluster with the given argument types. Values of this type are introduced as part of closure conversion.

T\_CFun **of** c\_prototype

the address of a C function with the given prototype.

T\_CStruct **of** ty

the address of memory containing a C struct value. This type is used to specify struct parameters and results in C function prototypes.

### 2.3 Kind W or V types

There are five numeric types whos representation (*i.e.*, kind) depends on the target architecture and compiler configuration. These types have either W kind, when they can be mapped into general-purpose registers or V kind when they cannot be so mapped. The types are:

#### T\_Int

32-bit 2's complement integers; we use int to denote this type.

#### T\_Long

64-bit 2's complement integers; we use **long** to denote this type.

#### T\_Float

32-bit IEEE single-precision floating-point numbers; we use float to denote this type.

T\_Double

64-bit IEEE double-precision floating-point numbers; we use **double** to denote this type.

T\_Extended

IEEE extended double-precision floating-point numbers; we use **extended** to denote this type.

## 2.4 Kind M types

T\_Data

a region of memory of unknown size.

T\_Object

a region of memory used to represent a MOBY object.

- T\_Struct of {sz : int, align : int, data : field list}
   a region of memory with a known size, alignment, and layout.
- T\_Vector of {len : int option, elemSz : int, ty : ty}
  an immutable vector of elements with the given size and type. When the len field is note
  NONE, then the length of the vector is known.

T\_Array of {len : int option, elemSz : int, ty : ty}
 a mutable array of elements with the given size and type. When the len field is note NONE,
 then the length of the vector is known.

T\_Union **of** ty list an untagged union of types.

T\_TaggedUnion **of** (int \* ty) list

### 2.5 Kind Type types

T\_Void

This type is used to denote the C void type in function prototypes.

# **3** Representation of Moby types

This section describes how common MOBY types are mapped to BOL types. It is always the case that the BOL type corresponding to a MOBY type will have W kind.

Bool is represented by the bool type.

Char is represented by (0..255).

Int is either represented by wrap(int) or by int.

Long is either represented by wrap(long) or by long.

Integer is represented by integer.

Float is represented by wrap(float).

Double is represented by wrap(double).

Extended is represented by wrap(extended).

#### **3.1** Sequence types

MOBY sequence types, such as Array and String, have a two-level representation in BOL. There is a two-word header consisting of a 32-bit integer length and a pointer to the data object.<sup>1</sup>

#### 3.2 The List type constructor

The MOBY List type constructor is defined as

```
datatype List(t) { Nil, Cons of (t, List(t)) }
```

The Nil value is represented by the value 0, while the Cons values are represented by pointers to two-word pairs. The BOL type for this representation is  $*\tau + (0..0)$ , where  $\tau$  is the type of the list elements (any when the type is unknown).

<sup>&</sup>lt;sup>1</sup>On 64-bit machines, there is 32-bits of padding between the length and the data pointer to ensure 64-bit alignment of the data pointer.

# 4 The BOL representation

Inside the MOBY compiler, BOL expressions are represented using the following datatypes:

```
datatype exp = E_Pt of (ProgPt.ppt * term)
and term = ...
and rhs = ...
```

The exp type is a term tagged with a unique program point. Program points serve as labels for those analyses that need to track positions in the code. The rhs (right-hand-side) type covers terms that cannot appear in a tail context.

#### 4.1 Expression forms

The term type has a number of constructors; we call these *expression forms* (ignoring the lack of a program-point label).

```
E_Let of (var list * exp * exp)
```

binds the variables to the results of the first expression in the scope of the second expression. The general syntax of this form is

let  $(x_1, ..., x_n) = e_1; e_2$ 

When the number of bound variables is one, we write

**let**  $x = e_1; e_2$ 

and when there are no bound variables, we write

do  $e_1$ ;  $e_2$ 

```
E_Stmt of (var list * rhs * exp)
```

binds the variables to the results of the right-hand-side in the scope of the expression. The syntax of this form is the same as for  $E\_Let$ .

```
E_StackAlloc of (var * int * int * exp)
```

binds the variable to reserved space in the stack frame. The first integer specifies the size (in bytes) of the space and the second specifies the alignment. The scope of the binding and the extent of the reserved space is the expression. The syntax for this form is

stackalloc  $x = \langle sz , align \rangle; e_2$ 

E\_Fun **of** (lambda list \* exp)

Binds a collection of mutually recursive function definitions. The scope of the function names includes both the function bodies and the expression. We use the syntax

```
fun f_1 (x_{1,1}, \ldots, x_{1,n_1}) = e_1
and \cdots
and f_k (x_{k,1}, \ldots, x_{k,n_k}) = e_k;
e
```

for the term

```
\texttt{E}\_\texttt{Fun}(\texttt{[}(f_1, \texttt{[}x_{1,1}, \ldots, x_{1,n_1}\texttt{]}, e_1), \ldots, \texttt{(}f_k, \texttt{[}x_{k,1}, \ldots, x_{k,n_k}\texttt{]}, e_k)\texttt{]}, e)
```

```
E_Cont of (lambda * exp)
```

Binds a BOL continuation with the expression as its scope. Note that the lifetime of the continuation is also its scope!

```
E_If of (var * exp * exp)
```

tests the variable and if it is true, the evaluate the first expression, otherwise evaluate the second expression. The syntax for this form is

if x then  $e_1$  else  $e_2$ 

E\_Switch of (var \* (int \* exp) list \* exp option)

Tests the variable against the integer labels of the list of cases; the third argument is the optional default case. The cases should be in increasing numeric order and the default case should be present unless the variable is guaranteed to always have one of the case labels as its values. We use the syntax

```
switch x { case i_1: e_1 \cdots case i_n: e_n }
```

for the term

```
E_Switch(x, [(i_1, e_1), ..., (i_n, e_n)], NONE)
```

and

switch  $x \{ \text{ case } i_1 : e_1 \cdots \text{ case } i_n : e_n \text{ default} : e \}$ 

for the term

```
\texttt{E\_Switch}(x, [(i_1, e_1), \ldots, (i_n, e_n)], \texttt{SOME}(e))
```

E\_Apply **of** (var \* var list)

applies the function named by the first variable to the arguments named by the list of variables. We use the syntax **call** f(args) for  $E_Apply(f, args)$ .

E\_Throw of (var \* var list)

applies the continuation named by the first variable to the arguments named by the list of variables. We use the syntax **throw** k(args) for E\_Throw(k, args).

```
E_Ret of var list
```

returns the values bound to the variables. Note that the term "*return*" does not connote control-flow.

#### 4.2 Lambda abstractions

The type lambda is used to represent both functions and continuations. It is defined as:

```
type lambda = (var * var list * exp)
```

where the first variable is the name of the function (there are no anonymous functions in BOL), the list of variables are the formal parameters, and the expression is the function body.

#### 4.3 Right-hand-side forms

```
E_Cast of (var * BOLTypes.ty)
```

cast the value bound to the variable to the given type (which must have the same kind). We use the notation  $(\tau)x$  for E\_Cast(x,  $\tau$ )

```
E_Select of (int * var)
```

selects the the specified field from the record bound to the variable. We use the notation x # i for E\_Select(*i*, *x*).

E\_Update of (int \* var \* var)

updates the specified field from the record bound to the first variable with the value bound to the second variable. This form has no results (*i.e.*, zero-arity). We use the notation x # i := y for E\_Update(*i*, *x*, *y*).

```
E_Alloc of (BOLTypes.ty * var list)
```

allocates and initialized a record in the heap. The type specifies the record's layout and the list of variables provide the initial values for record's fields.

```
E_AllocObj of (BOLTypes.ty * var)
```

allocate memory for an object. The type specifies the layout of the object's fields and the variable is bound to the method suite.

E\_Wrap **of** var

wrap (box) the value bound to the variable. We use the syntax wrap(x) for  $E_Wrap(x)$ .

E\_Unwrap **of** var

unwrap (unbox) the boxed value bound to the variable. We use the syntax unwrap(x) for E\_Unwrap(x).

E\_IConst **of** IntInf.int an integer constant.

E_SConst <b>of</b> string a string constant. Note that this is the string data and not the representation of a MOBY string literal.
E_FConst <b>of</b> FloatLit.float a floating-point constant.
E_BConst <b>of</b> bool a boolean constant.
E_StaticAddr <b>of</b> var the address of the static location named by the variable.
E_StaticRef <b>of</b> var the contents of the static location named by the variable.
E_Prim of var primop applies a primitive operator to its arguments. The primitive operators are described in Sec- tion 6.
E_Slot <b>of</b> slot_exp
E_DictFieldSel <b>of</b> (var * member_label)
E_DictMethSel <b>of</b> (var * member_label)
E_FieldGet <b>of</b> (var * var)
E_FieldPut <b>of</b> (var * var * var)
E_MethGet <b>of</b> (var * var)
E_ApplyCont of (var * var list) Partially apply a continuation to its arguments (but do not transfer control). This operation has the effect of turning a continuation with arguments into one without.

E\_ThdCreate of var

E\_ThdGetTask

#### 4.4 Creating BOL expressions

The BOL module provides constructor functions for the various expression forms (*e.g.*, mkLet to create an E\_Let expression form). These constructor functions take care of labeling the term with a unique program point. The Census module provides similar functions, except that they maintain the additional invariants defined by the census, such as variable binding information.

We use the syntax **ccall** f(args) for E\_CCall(f, args).

# 5 BOL variables

The representation of BOL variables has the SML type var, which is defined in the BOL module as follows:

```
datatype var = V of {
    id : Word.word,
    name : string option,
    src : Var.var option,
    binding : var_binding ref,
    ty : BOLTypes.ty,
    useCnt : int ref,
    props : PropList.holder
}
```

The fields of this representation are used as follows:

id a unique ID that can be used for identity testing, ordering, or hashing.

name if present, a symbolic name for the variable.

src if present, then this BOL variable corresponds to the specified typed AST variable.

binding the binding that defines this variable.

ty this variable's type.

useCnt the number of times that this variable is used. For functions and continuations, this count includes applications.

props a holder for name/value pairs (i.e., an association list).

## 6 **Primitive operators**

Machine-level operations are represented in BOL as "*primops*" (primitive operations). The primop datatype is defined in the PrimOps structure. This datatype is type constructor over the type used to represent the primop arguments; the BOL uses this type constructor applied to the var type. To ease the addition of new primitive operations, we generate the definition of the primop datatype and the various modules that directly work on it (*e.g.*, constant folding, effect analysis, code generation, *etc.*) from a specification file. The primitive operations can be grouped into the following classes:

- **Boolean operations** The boolean type serves as the result of conditionals and as the argument of conditionals. There is one operation logical negation.
- **Integer operations** There are two fixed-precision integer types in BOL: 32-bit and 64-bit. Each of these types has a complete set of arithmetic and comparison operations; the former are prefixed by "I32," while the latter are prefixed by "I64." In addition, there are unsigned comparisons on 32-bit integers (prefixed by "U32").

- **Floating-point operations** There are three floating-point types: IEEE 32-bit single-precision numbers, IEEE 64-bit double-precision numbers, and IEEE extended-double-precision numbers. The size of the latter type depends on the target architecture; it is 80-bits on the Intel IA32 (a.k.a. x86) and 64-bits on the PowerPC. Each of these types has a complete set of arithmetic and comparison operations that follow the IEEE semantics. In addition, there are two multiply accumulate instructions that can produce non-IEEE results.
- **String operations** BOL provides operations for comparison of string data values. Since these values do not have length information (see Section 3), they take a first argument which is a limit on the number of characters to compare.
- **Pointer testing operations** The translation of higher-level datatypes (*e.g.*, lists) uses the distinction between pointers and small integers (integers in the range  $[0, 2^{16} 1]$ ) to distinguish between different constructors. In this case, we call the pointer a *boxed* value and the small integer a *unboxed* value. BOL provides operations to test for boxed and unboxed values.
- Address arithmetic BOL has a full complement of address arithmetic operations. These are used to support data-level interoperability with foreign code and data structures.
- **Conversion operations** BOL has conversion operators between the various numeric types. In addition, it has operations to cast between integer and floating-point representations (*e.g.*, to allow one to examine the bits of a floating-point number directly.
- Synchronization operations BOL includes low-level synchronization operations to support spin locks and the like.

The following is a list of the BOL primitive operations with their types and a short description of each operator:

BNot	: Bool -> Bool Boolean negation.
I32Neg	: Int -> Int
	32-bit 2's complement negation.
I32Add	: (Int, Int) -> Int
	32-bit 2's complement addition.
I32Sub	: (Int, Int) -> Int
	32-bit 2's complement subtraction.
I32Mul	: (Int, Int) -> Int
	32-bit 2's complement multiplication.
I32Div	: (Int, Int) -> Int
	32-bit 2's complement division.
I32Mod	: (Int, Int) -> Int
	32-bit 2's complement remainder.

I32Not	:	Int -> Int 32-bit 1's complement negation
I32And	:	(Int, Int) -> Int
		32-bit logical and.
1320r	:	(Int, Int) -> Int 32-bit logical or.
I32XOr	:	(Int, Int) -> Int
		32-bit logical xor.
I32LSh	:	(Int, Int) -> Int 32-bit left-shift
T32RSha		$(\text{Int} \text{Int}) \rightarrow \text{Int}$
1 J Z KOIIA	•	32-bit arithmetic right-shift
T32RShL	:	$(\text{Int}, \text{Int}) \rightarrow \text{Int}$
		32-bit logical right-shift.
T 3 2 T.+	:	$(Int, Int) \rightarrow Bool$
19210		32-bit 2's complement less-than comparison
T32Lte	:	(Int. Int.) -> Bool
101100		32-bit 2's complement less-than or equal comparison.
I32Gt	:	(Int. Int) -> Bool
10100		32-bit 2's complement greater comparison.
T32Gte	:	(Int. Int.) -> Bool
101000		32-bit 2's complement greater-than or equal comparison.
T64Ea	:	(Int. Int.) -> Bool
1		64-bit equal test.
I64NEa	:	(Int, Int) -> Bool
1		64-bit not-equal test.
		1
U32Lt	:	(Int, Int) -> Bool
		32-bit unsigned less-than comparison.
U32Lte	:	(Int, Int) -> Bool
		32-bit unsigned less-than or equal comparison.
U32Gt	:	(Int, Int) -> Bool
		32-bit unsigned greater comparison.
U32Gte	:	(Int, Int) -> Bool
		32-bit unsigned greater-than or equal comparison.
I64Neg	:	Int -> Int
		64-bit 2's complement negation.
I64Add	:	(Int, Int) -> Int
		64-bit 2's complement addition.
I64Sub	:	(Int, Int) -> Int
		64-bit 2's complement subtraction.
I64Mul	:	(Int, Int) -> Int
		64-bit 2's complement multiplication.
I64Div	:	(Int, Int) -> Int

		64-bit 2's complement division.
T64Mod	:	(Int. Int.) -> Int
1011100		64-bit 2's complement remainder.
T64Not	:	Int -> Int
10 11000		64-bit 1's complement negation
T64And		$(Tnt Tnt) \rightarrow Tnt$
TOTAIL	•	64-bit logical and
T610r		$(1 \text{ pt} 1 \text{ pt}) = \sum_{i=1}^{n} 1 \text{ pt}$
10401	•	$(110, 110) \rightarrow 110$
TGAYOr		$(\text{Trt}, \text{Trt}) = \sum_{n=1}^{n} \text{Trt}$
104X01	•	$(110, 110) \rightarrow 110$
TEATCH		$(1 \text{ nt} 1 \text{ nt}) = \sum_{i=1}^{n} 1 \text{ nt}$
1040311	•	$(110, 110) \rightarrow 110$
TEADCHA		$(1 \text{ Trt} \text{ Trt}) = \sum_{i=1}^{n} 1 \text{ Trt}$
TOARDIA	•	(IIIC, IIIC) -> IIIC
TEADCHT		$(1 \text{ trat} \text{ trat}) = \sum_{i=1}^{n} 1 \text{ trat}$
104RSIIL	•	(IIIC, IIIC) -> IIIC 64 bit logical right shift
тблтн		$(1 \text{ tr} t + 1 \text{ tr} t) = \sum_{n=1}^{\infty} P_{n} P_{n}$
	•	64 bit 2's complement less than comparison
TEATEO		$(1 \text{ pt} - 1 \text{ pt}) = \sum_{i=1}^{n} \sum_{j=1}^{n} \sum_{i$
TOATICE	•	64 bit 2's complement less than or equal comparison
T64C+		(Int Int) -> Bool
10490	•	64-bit 2's complement greater comparison
T64Cto	•	(Int Int) -> Bool
101000	•	64-bit 2's complement greater-than or equal comparison
T64Ea	:	(Int Int) -> Bool
10104		64-bit equal test
T64NEC	:	$(\text{Int} \text{Int}) \rightarrow \text{Bool}$
10 11104		64-bit not-equal test
		or on not equal test.
F32Neg	:	Float -> Float
		32-bit IEEE floating-point negation
F32Add	:	(Float, Float) -> Float
		32-bit IEEE floating-point addition
F32Sub	:	(Float, Float) -> Float
		32-bit IEEE floating-point subtraction
F32Mul	:	(Float, Float) -> Float
		32-bit IEEE floating-point multiplication
F32Div	:	(Float, Float) -> Float
		32-bit IEEE floating-point division
F32Rem	:	(Float, Float) -> Float
		32-bit IEEE floating-point remainder
F32MAdd	:	(Float, Float, Float) -> Float
		32-bit floating-point multiply/add
F32MSub	:	(Float, Float, Float) -> Float

		32-bit floating-point multiply/subtract
F32Abs	:	Float -> Float
		32-bit IEEE floating-point absolute value
F32CopySign	:	(Float, Float) -> Float
192002791911		32-bit IEEE floating-point copy-sign
F32Sart		Float -> Float
FJZBYLU	•	22 hit IEEE floating point square root
<b>H</b> 20D		(Plast Plast) > Plast
F3ZPOW	•	(Float, Float) -> Float
F32Lt	:	(Float, Float) -> Bool
		32-bit IEEE floating-point less-than comparison.
F32Lte	:	(Float, Float) -> Bool
		32-bit IEEE floating-point less-than or equal comparison.
F32Gt	:	(Float, Float) -> Bool
		32-bit IEEE floating-point greater-than comparison.
F32Gte	:	(Float, Float) -> Bool
		32-bit IEEE floating-point greater-than or equal comparison.
F32Eq	:	(Float, Float) -> Bool
1		32-bit IEEE floating-point inequality test
F32NFC	•	(Float Float) => Bool
roznag	•	32 bit IEEE floating point equality test
		(Elect Elect) -> Peel
FSZLLGL	•	(Float, Float) -> Bool
FSZULL	•	(FIOAL, FIOAL) -> BOOI
F3ZULLE	•	(Float, Float) -> Bool
F32UGt	:	(Float, Float) -> Bool
F32UGte	:	(Float, Float) -> Bool
F320rdered	:	(Float, Float) -> Bool
		32-bit IEEE floating-point ordered test.
F32Unordered	:	(Float, Float) -> Bool
		32-bit IEEE floating-point unordered test.
F32Finite	:	Float -> Bool
		test for 32-bit IEEE finite number
F32Infinite	:	Float -> Bool
		test for 32-bit IEEE infinite number
F64Neg		Double -> Double
1 0 11109	•	64-bit IEEE floating-point negation
E64744		$(\text{Double}  \text{Double}) = \sum \text{Double}$
I UTAUU	•	64 bit IEEE floating point addition
F64SUD	:	(Double, Double) -> Double

		64-bit IFFE floating-point subtraction
F64Mul		$(\text{Double}  \text{Double}) = \sum \text{Double}$
r o finar	•	64-bit IEEE floating-point multiplication
		$(\text{Devbl}_{2}, \text{Devbl}_{2}) = \sum_{i=1}^{n} \text{Devbl}_{2}$
F04DIV	•	(Double, Double) -> Double
		64-bit IEEE floating-point division
F64Rem	:	(Double, Double) -> Double
		64-bit IEEE floating-point remainder
F64MAdd	:	(Double, Double, Double) -> Double
		64-bit floating-point multiply/add
F64MSub	:	(Double, Double, Double) -> Double
		64-bit floating-point multiply/subtract
F64Abs	:	Double -> Double
		64-bit IEEE floating-point absolute value
F64CopvSign	:	(Double, Double) -> Double
		64-bit IEEE floating-point copy-sign
F64Sart	:	Double -> Double
roibqre	•	64-bit IFFF floating-point square root
E61Dow		(Double Double) -> Double
FOFFOW	•	(Double, Double) -> Double
		(Deuble Deuble) > Deel
F04LL	•	(Double, Double) -> Bool
		64-bit IEEE floating-point less-than comparison.
F64Lte	:	(Double, Double) -> Bool
		64-bit IEEE floating-point less-than or equal comparison.
F64Gt	:	(Double, Double) -> Bool
		64-bit IEEE floating-point greater-than comparison.
F64Gte	:	(Double, Double) -> Bool
		64-bit IEEE floating-point greater-than or equal comparison.
F64Eq	:	(Double, Double) -> Bool
		64-bit IEEE floating-point inequality test.
F64NEq	:	(Double, Double) -> Bool
		64-bit IEEE floating-point equality test.
F64LtGt	:	(Double, Double) -> Bool
F64III.t	:	$(Double, Double) \rightarrow Bool$
1010110		(Double, Double, P Dool
		$(\text{Double}, \text{Double}) \rightarrow \text{Real}$
rofonce	•	(Double, Double) -> Bool
F64UGC	•	(Double, Double) -> Bool
F64UGte	:	(Double, Double) -> Bool
F640rdered	:	(Double, Double) -> Bool
		64-bit IEEE floating-point ordered test.
F64Unordered	:	(Double, Double) -> Bool
		64-bit IEEE floating-point unordered test.

F64Finite	:	Double -> Bool test for 64-bit IEEE finite number
F64Infinite	:	Double -> Bool test for 64-bit IEEE infinite number
FXNeg	:	Extended -> Extended
FXAdd	:	(Extended, Extended) -> Extended
FXSub	:	(Extended, Extended) -> Extended
FXMul	:	(Extended, Extended) -> Extended
FXDiv	:	(Extended, Extended) -> Extended
FXRem	:	(Extended, Extended) -> Extended
FXMAdd	:	(Extended, Extended, Extended) -> Extended
FXMSub	:	(Extended, Extended, Extended) -> Extended
FXAbs	:	Extended -> Extended
FXCopySign	:	(Extended, Extended) -> Extended
FXSqrt	:	Extended -> Extended
FXPow	:	(Extended, Extended) -> Extended
FXLt	:	(Extended, Extended) -> Bool
FXLte	:	(Extended, Extended) -> Bool
FXGt	:	(Extended, Extended) -> Bool
FXGte	:	(Extended, Extended) -> Bool
FXEq	:	(Extended, Extended) -> Bool
FXNEq	:	(Extended, Extended) -> Bool
FXLtGt	:	(Extended, Extended) -> Bool
FXULt	:	(Extended, Extended) -> Bool

FXULte	:	(Extended, Extended) -> Bool
FXUGt	:	(Extended, Extended) -> Bool
FXUGte	:	(Extended, Extended) -> Bool
FXOrdered	:	(Extended, Extended) -> Bool
FXUnordered	:	(Extended, Extended) -> Bool
FXFinite	:	Extended -> Bool
FXInfinite	:	Extended -> Bool
StrEq	:	(Int, String, String) -> Bool test two strings for equality
StrNEq	:	(Int, String, String) -> Bool test two strings for inequality
StrCmp	:	(Int, String, String) -> Int compare two strings for order
Boxed	:	Any -> Bool test for boxed values
Unboxed	:	Any -> Bool test for unboxed values
AdrEq	:	(Addr, Addr) -> Bool
AdrNEq	:	(Addr, Addr) -> Bool test addresses for inequality
AdrAdd	:	(Addr, Int) -> Addr add an integer to an address
AdrSub	:	(Addr, Int) -> Addr subtract an integer from an address
AdrAdd4	:	(Addr, Int) -> Addr add a scaled (by 4) integer to an address
AdrSub4	:	(Addr, Int) -> Addr
AdrAdd8	:	(Addr, Int) -> Addr
AdrSub8	:	(Addr, Int) -> Addr
AdrLoadI8	:	Addr -> Int load a sign-extended 8-bit integer from memory

AdrStoreI8	:	(Addr, Int) -> () store an 8-bit integer
Ndrt ord T16		Addr Int
Adilloadilb	•	load a sign-extended 16-bit integer from memory
AdrStoreI16	:	(Addr. Int) -> ()
		store a 16-bit integer
AdrLoadI32	:	Addr -> Int
		load a 32-bit integer from memory
AdrStoreI32	:	(Addr, Int) -> ()
		store a 32-bit integer
AdrLoadI64	:	Addr -> Long
		load a 64-bit integer from memory
AdrStoreI64	:	(Addr, Long) -> ()
		store a 64-bit integer
AdrLoadF32	:	Addr -> Float
		load a 32-bit floating-point number from memory
AdrStoreF32	:	(Addr, Float) -> ()
		store a 32-bit floating-point number
AdrLoadF64	:	Addr -> Double
		load a 64-bit floating-point number from memory
AdrStoreF64	:	(Addr, Double) -> ()
		store a 64-bit floating-point number
AdrLoadFX	:	Addr -> Extended
		load an extended-precision floating-point number from memory
AdrStoreFX	:	(Addr, Extended) -> ()
		store a extended-precision floating-point number
AdrLoadP	:	Addr -> Addr
		load an address from memory
AdrStoreP	:	(Addr, Ptr) -> ()
		store an address
AdrLoadU8	:	Addr -> Int
		load an unsigned 8-bit integer from memory
AdrLoadU16	:	Addr -> Int
		load an unsigned 16-bit integer from memory
AdrLoad	:	Addr -> Any
		load a word from memory
AdrStore	:	(Addr, Any) -> ()
		store a word
CVT13210164	÷	LILL -> LONG
Q		zero-extend a 52-bit integer to a 64-bit integer.
CVTXI32TO164	÷	LILL -> LONG
		sign-extend a 52-bit integer to a 64-bit integer.
CVT13210F32	:	Int -> Float
		convert a 32-bit integer to a 32-bit floating-point number.

CvtI32ToF64	:	Int -> Double
		convert a 32-bit integer to a 64-bit floating-point number.
CvtI32ToFX	:	Int -> Extended
		convert a 32-bit integer to an extended-precision floating-point number.
CastF32ToI32	:	Int -> Float
Cast I32ToE32		The -> Float
Cascisziorsz	•	cast a 32-bit integer to a 32-bit floating-point number.
CastF64ToI64	:	Double -> Long
		cast a 64-bit floating-point number to a 64-bit integer.
CastI64ToF64	:	Double -> Long
		cast a 64-bit integer to a 64-bit floating-point number.
CvtF32ToF64	:	Float -> Double
		convert a 32-bit floating-point number to a 64-bit floating-point number.
CvtF32ToFX	:	Float -> Extended
		convert a 32-bit floating-point number to an extended-precision floating-point number.
CvtF64ToFX	:	Double -> Extended
		convert a 64-bit floating-point number to an extended-precision floating-point number.
I32CmpAndSwap	:	(Addr, Int, Int) -> Bool, Int
I64CmpAndSwap	:	(Addr, Long, Long) -> Bool, Long