

## CS281 Spring 2010: Homework 3

Due Wednesday, April 28th – in class

1. (10 points)

Informally but clearly describe a Nondeterministic Turing machine – multitape if you like – that accepts the following language [Try to take advantage of nondeterminism to avoid iteration and save time in the nondeterministic sense. That is, prefer to have your NTM branch a lot, while each branch is short.] :

The language of all strings of the form  $w_1\#w_2\#\dots\#w_n$ , for any  $n$ , such that each  $w_i$  is a string of 0's and 1's, and for some  $j$ ,  $w_j$  is the integer  $j$  in binary.

2. (10 points)

Consider the nondeterministic Turing machine

$$M = (\{q_0, q_1, q_2, q_f\}, \{0, 1\}, \{0, 1, B\}, \delta, q_0, B, \{q_f\})$$

Informally but clearly describe the language  $L(M)$  if  $\delta$  consists of the following sets of rules:  $\delta(q_0, 0) = \{(q_0, 1, R), (q_1, 1, R)\}$ ;  $\delta(q_1, 1) = \{(q_2, 0, L)\}$ ;  $\delta(q_2, 1) = \{(q_0, 1, R)\}$ ;  $\delta(q_1, B) = \{(q_f, B, R)\}$ .

3. (10 points)

A  $k$ -head Turing Machine has  $k$  heads reading cells of one tape. A move of this TM depends on the state and on the symbol scanned by each head. In one move, the TM can change state, write a new symbol on the cell scanned by each head. and can move each head left, right or keep it stationary. Since several heads may be scanning the same cell, we assume the heads are numbered 1 through  $k$ , and the symbol written by the highest numbered head scanning a given cell is the one that actually gets written there. Prove that the languages accepted by  $k$ -head Turing Machines are the same as those accepted by ordinary TM's.

4. (10+10 = 20 points)

State (with justification) whether the recursive languages and the RE languages are closed under the following operations. You may give informal but clear constructions to show closure.

- Concatenation
- Kleene Star<sup>1</sup> operation.

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<sup>1</sup>Given a language  $L$ , the Kleene star of  $L$  is the language  $L^* = \bigcup_{n \in \mathbb{N}} L^n$  where  $\forall n > 0$ ,  $L^n$  is the language consisting of concatenations of  $n$  elements of  $L$  and  $L^0 = \{\Lambda\}$ ,  $\Lambda$  being the empty string.