24. Hardware Security (Meltdown, Spectre, TEE) & Authentication Part 3

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Hardware Security: A Broad View

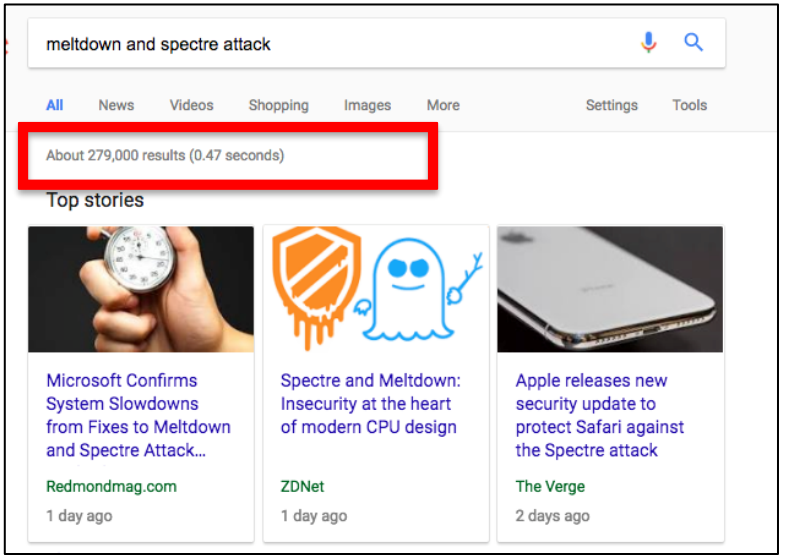
- What do we trust?
- How do we know we have the right code?
  - Recall software checksums, Subresource Integrity (SRI)
- What is our root of trust? Can we establish a smaller one?
- Can we minimize the Trusted Computing Base (TCB)?
- Can processor design lead to insecurity?
  - Yes! ☹

Attacks that exploit processor vulnerabilities

Can leak sensitive data

Relatively hard to mitigate

Lots of media attention



Relevant Ideas in CPUs

- **Memory isolation**: Processes should only be able to read their own memory
  - Virtual (paged) memory
  - Protected memory / Protection domains
- CPUs have a relatively small, and very fast, cache
  - Loading uncached data can take >100 CPU cycles
- **Out-of-order execution**: Order of processing in CPU can differ from the order in code
  - Instructions are much faster than memory access; you might be waiting for operands to be read from memory
  - Instructions **retire** (return to the system) in order even if they executed out of order

Relevant Ideas in CPUs

- There might be a conditional branch in the instructions
- **Speculative execution**: Rather than waiting to determine which branch of a conditional to take, go ahead anyway
  - **Predictive execution**: Guess which branch to take
  - **Eager execution**: Take both branches
- When the CPU realizes that the branch was mis-speculatively executed, it tries to eliminate the effects
- A core idea underlying Spectre/Meltdown: The results of the instruction(s) that were mis-speculatively executed will be cached in the CPU *[yikes!]*

Example (Not bad)

Consider the code sample below. If `arr1->length` is uncached, the processor can speculatively load data from `arr1->data[untrusted_offset_from_caller]`. This is an out-of-bounds read. That should not matter because the processor will effectively roll back the execution state when the branch has executed; none of the speculatively executed instructions will retire (e.g. cause registers etc. to be affected).

```
struct array {
 unsigned long length;
 unsigned char data[];
};
struct array *arr1 = ...;
unsigned long untrusted_offset_from_caller = ...;
if (untrusted_offset_from_caller < arr1->length) {
 unsigned char value = arr1->data[untrusted_offset_from_caller];
 ...
}
```

https://googleprojectzero.blogspot.com/2018/01/reading-privileged-memory-with-side.html

Example (Bad!!!)

However, in the following code sample, there's an issue. If `arr1->length`, `arr2->data[0x200]` and `arr2->data[0x300]` are not cached, but all other accessed data is, and the branch conditions are predicted as true, the processor can do the following speculatively before `arr1->length` has been loaded and the execution is re-steered:

- load `value = arr1->data[untrusted_offset_from_caller]`
- start a load from a data-dependent offset in `arr2->data`, loading the corresponding cache line into the L1 cache

```
struct array {
 unsigned long length;
 unsigned char data[];
};
struct array *arr1 = ...; /* small array */
struct array *arr2 = ...; /* array of size 0x400 */
/* >0x400 (OUT OF BOUNDS!) */
unsigned long untrusted_offset_from_caller = ...;
if (untrusted_offset_from_caller < arr1->length) {
 unsigned char value = arr1->data[untrusted_offset_from_caller];
 unsigned long index2 = ((value&1)*0x100)+0x200;
 if (index2 < arr2->length) {
   unsigned char value2 = arr2->data[index2];
 }
}
```

After the execution has been returned to the non-speculative path because the processor has noticed that `untrusted_offset_from_caller` is bigger than `arr1->length`, the cache line containing `arr2->data[index2]` stays in the L1 cache. By measuring the time required to load `arr2->data[0x200]` and `arr2->data[0x300]`, an attacker can then determine whether the value of `index2` during speculative execution was 0x200 or 0x300 - which discloses whether `arr1->data[untrusted_offset_from_caller]&1` is 0 or 1.

https://googleprojectzero.blogspot.com/2018/01/reading-privileged-memory-with-side.html

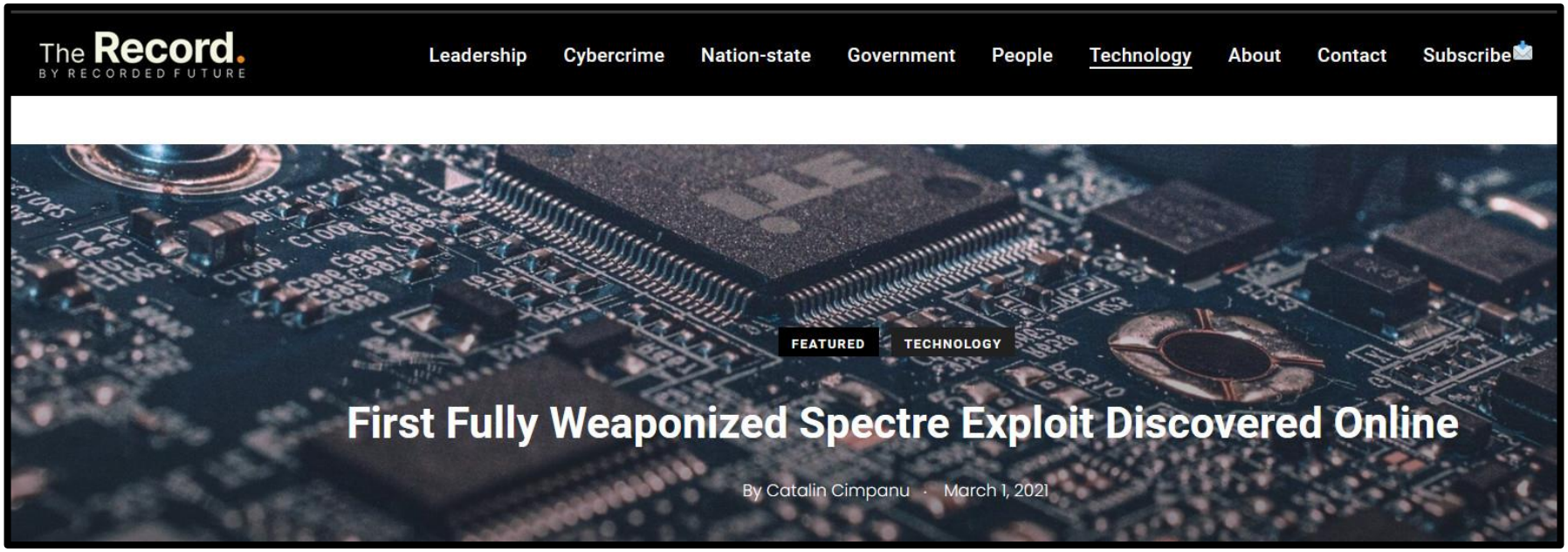
Spectre: Key Idea

- Use branch prediction as on the previous slide
- Conducting a timing side-channel attack on the cache
- Determine the value of interest based on the speed with which it returns
- Spectre allows you to read any memory from your process **for nearly every CPU**

Spectre: Exploitation Scenarios

- Leaking browser memory
- JavaScript (e.g., in an ad) can run Spectre
- Can leak browser cache, session key, other site data

Spectre: Exploitation Scenarios



"But today, Voisin said he discovered new Spectre exploits—one for Windows and one for Linux—different from the ones before. In particular, Voisin said he found a Linux Spectre exploit capable of dumping the contents of ***/etc/shadow***, a Linux file that stores details on OS user accounts"

https://therecord.media/first-fully-weaponized-spectre-exploit-discovered-online/

Meltdown: Key Idea

1. Attempt instruction with memory operand (Base+A), where A is a value forbidden to the process
2. The CPU schedules a privilege check and the actual access
3. The privilege check fails, but due to speculative executive, the access has already run and the result has been cached
4. Conduct a timing attack reading memory at the address (Base+A) for all possible values of A. The one that ran will return faster

Meltdown allows you to read **any memory in the address space (even from other processes)** but only on some Intel/ARM CPUs

Meltdown Attack (Timing)

- Now the attacker read each page of probe array
- 255 of them will be slow
- The X^{th} page will be faster (it is cached!)
- We get the value of X using cache-timing side channel

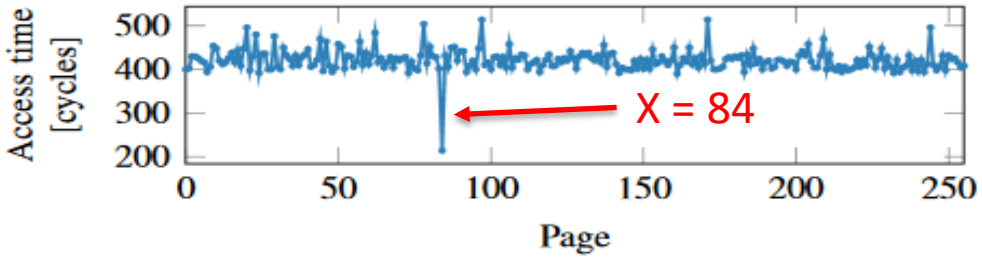


Figure 4: Even if a memory location is only accessed during out-of-order execution, it remains cached. Iterating over the 256 pages of `probe_array` shows one cache hit, exactly on the page that was accessed during the out-of-order execution.

Meltdown: Mitigation

- KAISER/KPTI (kernel page table isolation)
- Remove kernel memory mapping in user space processes
- Has non-negligible performance impact
- Some kernel memory still needs to be mapped

Trusted Computing

Trusted Platform Module (TPM)

- Standardization of cryptoprocessors, or microcontrollers dedicated to crypto functions w/ built-in keys
- Core functionality:

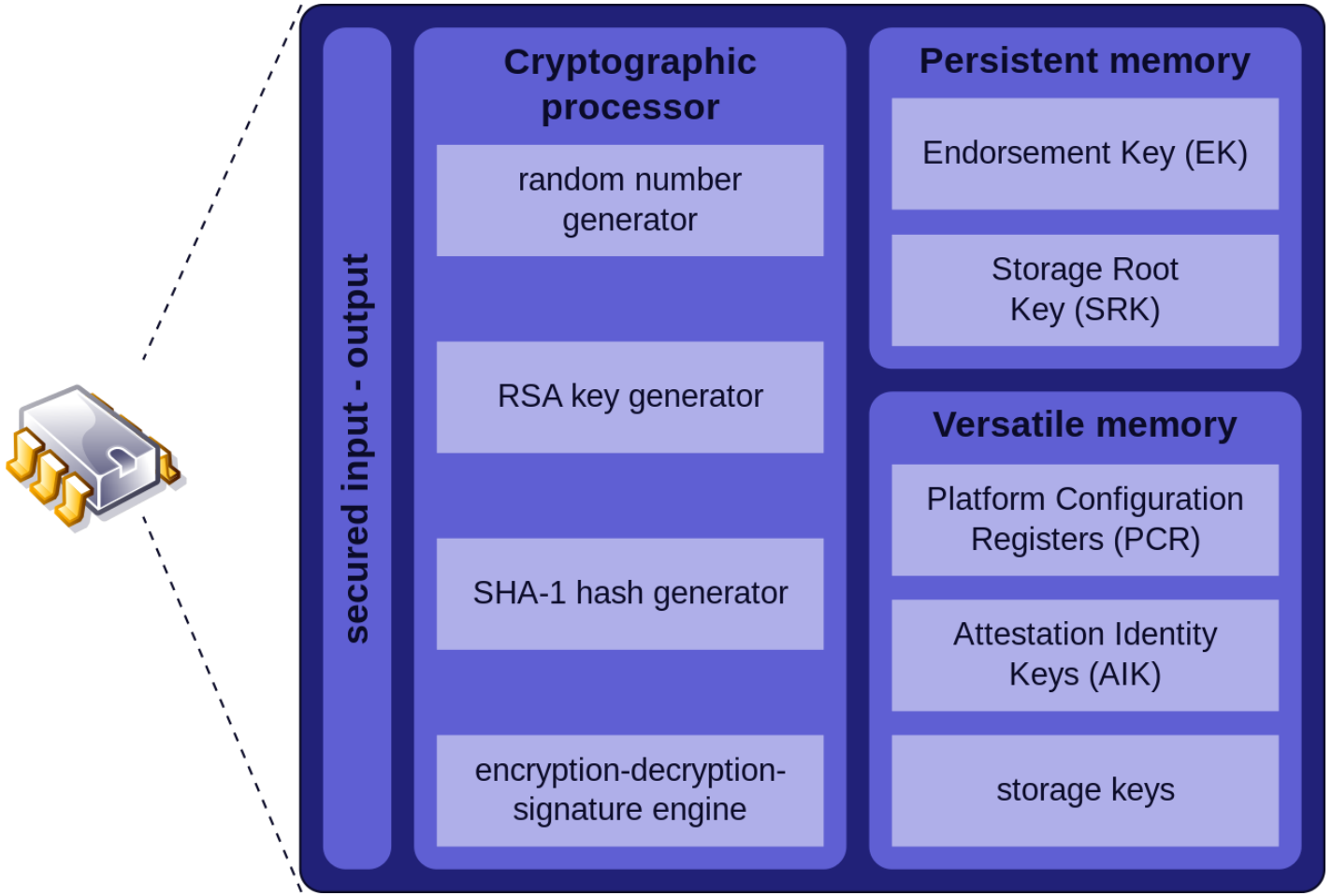
1) Random number generation, crypto key creation

2) **Remote attestation** (hash hardware and software config and send it to a verifier)

3) **Bind/seal** data: encrypted using a TPM key and, for sealing, also the required TPM state for decryption

- Uses: DRM, disk encryption (BitLocker), auth

Trusted Platform Module (TPM)

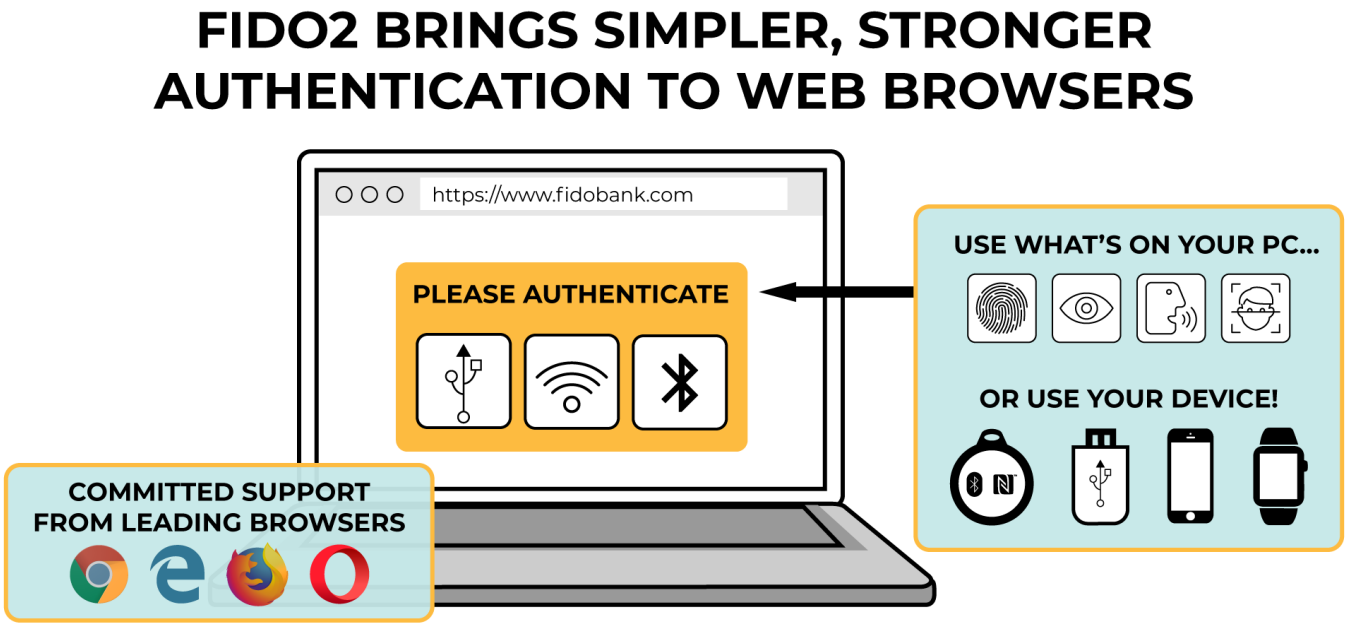


Trusted Execution Environment (TEE)

- TPMs are standalone companion chips, while TEEs are a secure area of a main processor
- Guarantees confidentiality and integrity for code in TEE
- Key example: Intel Software Guard Extensions (SGX)
- ***Enclaves* =** Private regions of memory that can't be read by any process outside the enclave, even with root access
- Uses: DRM, mobile wallets, auth

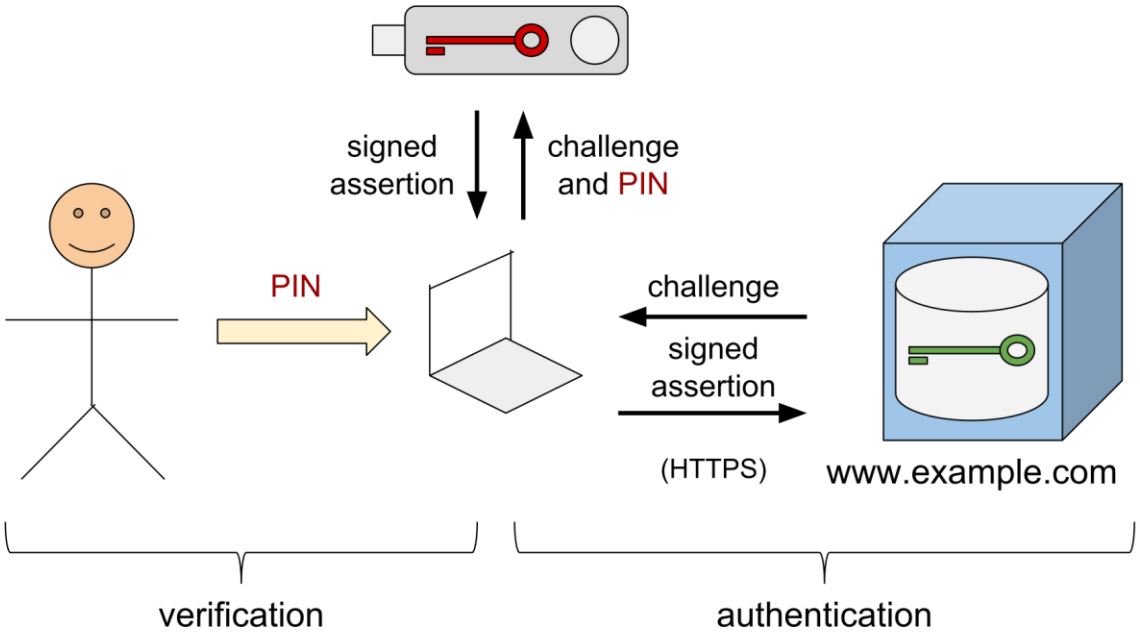
Authentication in Practice: Moving Towards A Passwordless World?

Case Study: WebAuthn



Case Study: WebAuthn

- Created under the FIDO2 project, now a W3C standard
- Goal: Authenticate on web using public-key crypto
- Implemented in specialized hardware OR in software using a TPM/TEE



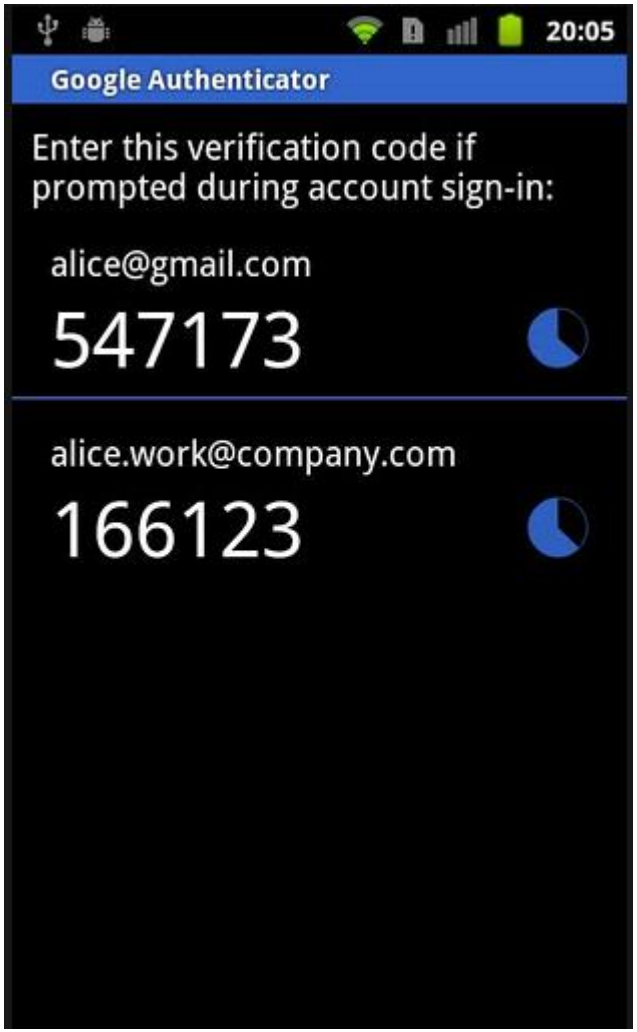
Case Study: WebAuthn

User interaction: Push a button on a key, type a PIN into the device, present biometric (fingerprint) to hardware reader

Authentication in Practice: Password Add-Ons / Alternatives

Single Sign-On

Two-Factor Auth



Physical Tokens / Smart Cards

- Codes based on a cryptographic key
  - Token manufacturer also knows the key
- What if there is a breach?

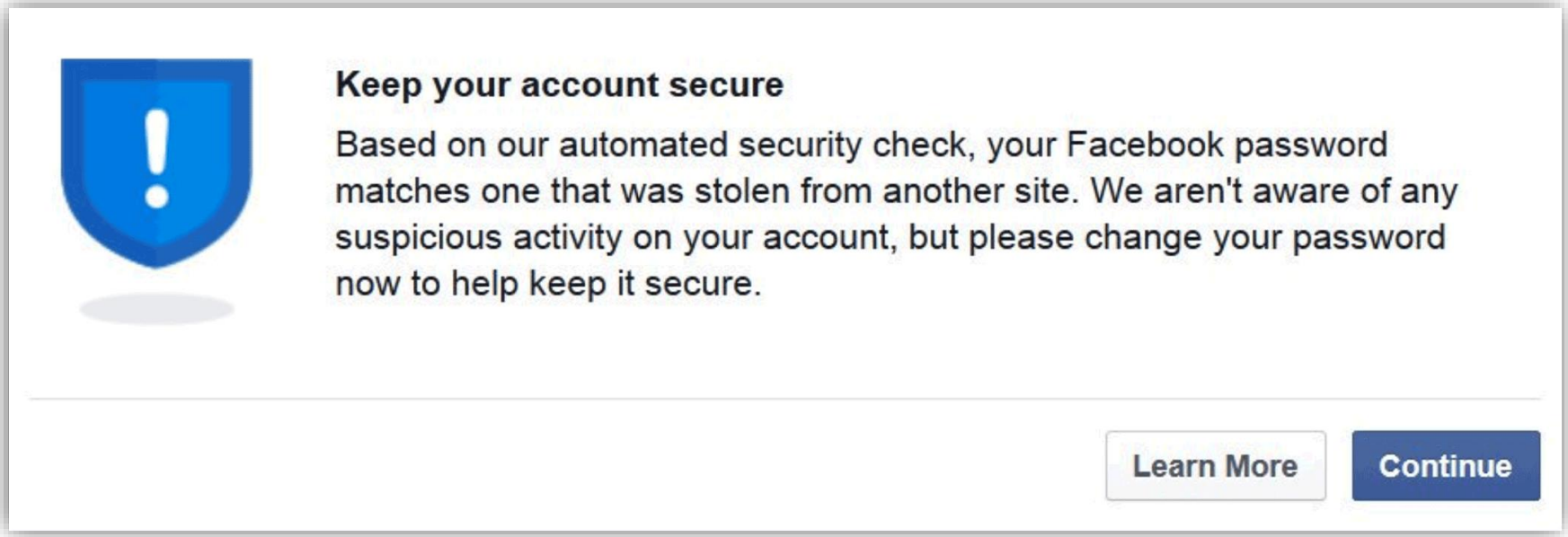
Authentication in Practice: I Forgot My Password

Resetting Accounts

- I forgot my password!
- Send an email?
- Security questions?
- In-person verification?
- Other steps?
- (No backup)

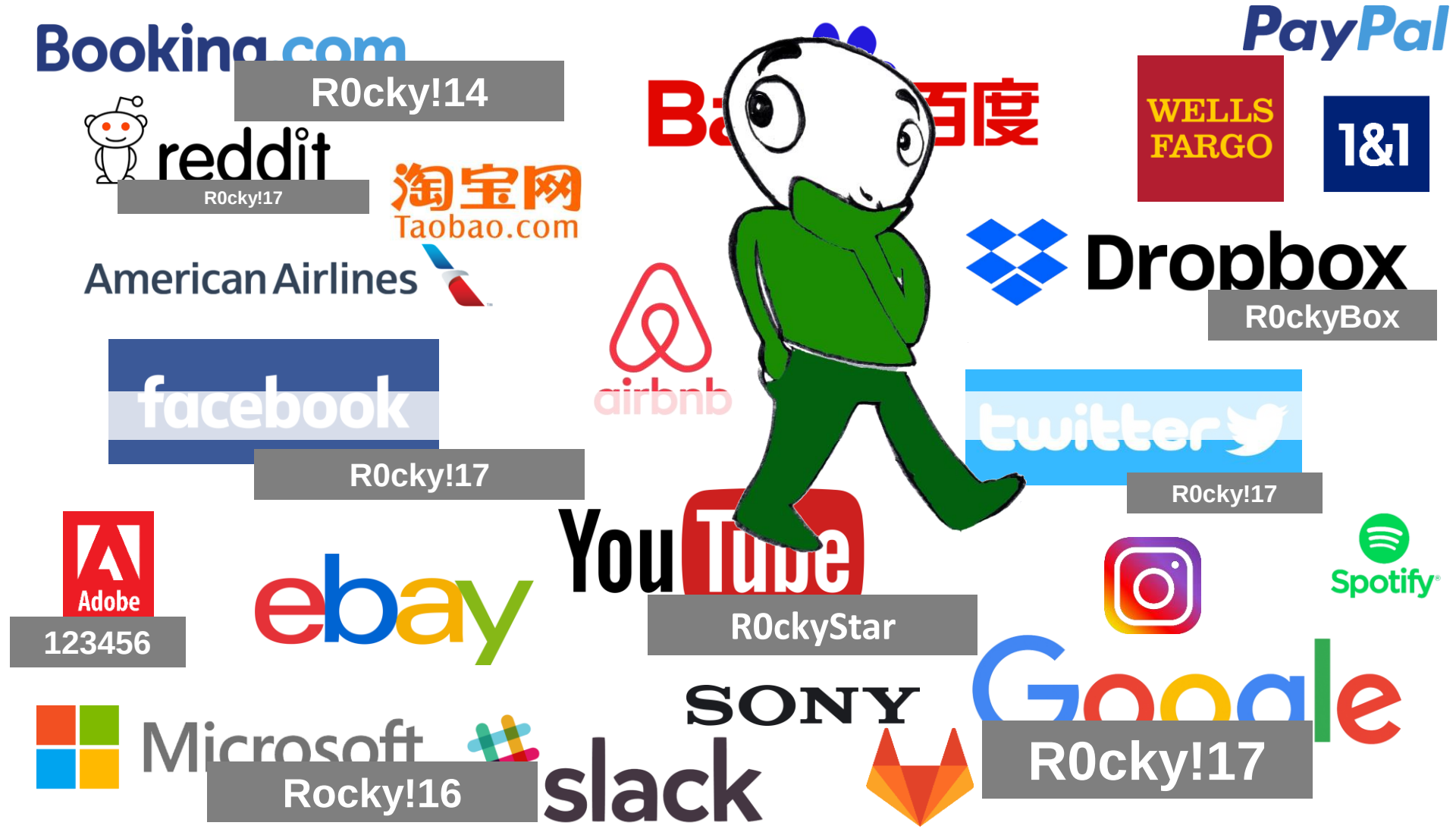
Authentication in Practice: Password Reuse ☹

Password Reuse-Based Attacks



Maximilian Golla, Miranda Wei, Juliette Hainline, Lydia Filipe, Markus Dürmuth, Elissa Redmiles, Blase Ur. "What was that site doing with my Facebook Password?" Designing Password-Reuse Notifications. In *Proc. CCS*, 2018.

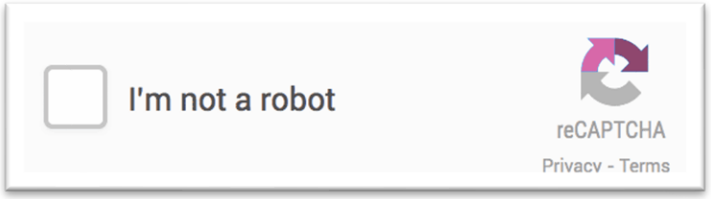
People Reuse Passwords



Memory-Hard Hash Function

Email	Argon2i Hash of Password
…	…
jim@mail.com	$argon2i$v=19$m=4096,…
…	…

Rate-Limiting Guessing



Password Strength Meter

Username

Password

acmccs18

Show Password & Detailed Feedback

Your password could be better.

- Consider inserting digits into the middle, not just at the end (Why?)
- Make your password longer than 8 characters (Why?)
- Consider using 1 or more symbols (Why?)

A better choice: \a#D18cmccs

How to make strong passwords


AcmeCo
Email
...
jim@mail.com
...


LinkedIn
Email
jane@aol.com
jessey@gmx.net
jenny@gmail.com
jim@mail.com
john@hotmail.com
...

LinkedIn

Email	SHA-1 Hash of Password
jane@aol.com	7c4a8d09ca3762af61e595209
jessey@gmx.net	5baa61e4c9b93f3f0682250b6
jenny@gmail.com	7c222fb2927d828af22f59213
jim@mail.com	ba93664a90285b9ff18a7a081
john@hotmail.com	b1b3773a05c0ed0176787a4f1
...	...

Crack All The Things!

```
Bash
$> hashcat -m 100 -a0 $TARGET $DICT
123456
Password
R0cky!17
Football!17
CanadaRocks!
```

LinkedIn

Email	Cracked SHA-1 Hashes
jane@aol.com	123456
jessey@gmx.net	5baa61e4c9b93f3f0682250b6
jenny@gmail.com	Canada4ever
jim@mail.com	R0cky!17
john@hotmail.com	HikingGuy89
…	…

Dead On Arrival

AcmeCo

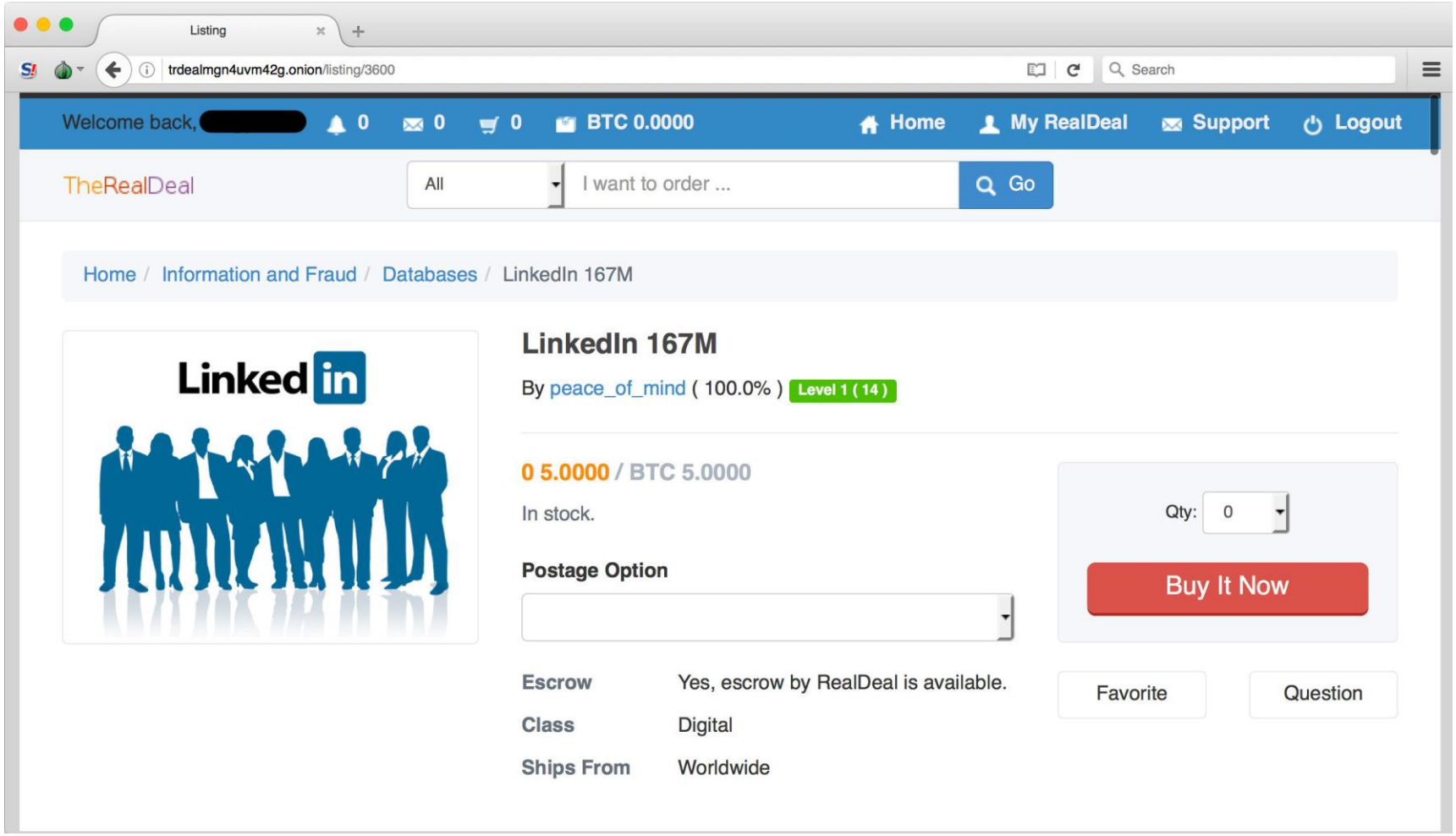
Email	Cracked
…	…
jim@mail.com	R0cky!17
…	…

1 guess is enough!

LinkedIn

Email	Cracked SHA-1 Hashes
jane@aol.com	123456
jessey@gmx.net	5baa61e4c9b93f3f0682250b6
jenny@gmail.com	Canada4ever
jim@mail.com	R0cky!17
john@hotmail.com	HikingGuy89
…	…

Monitoring the Black Market



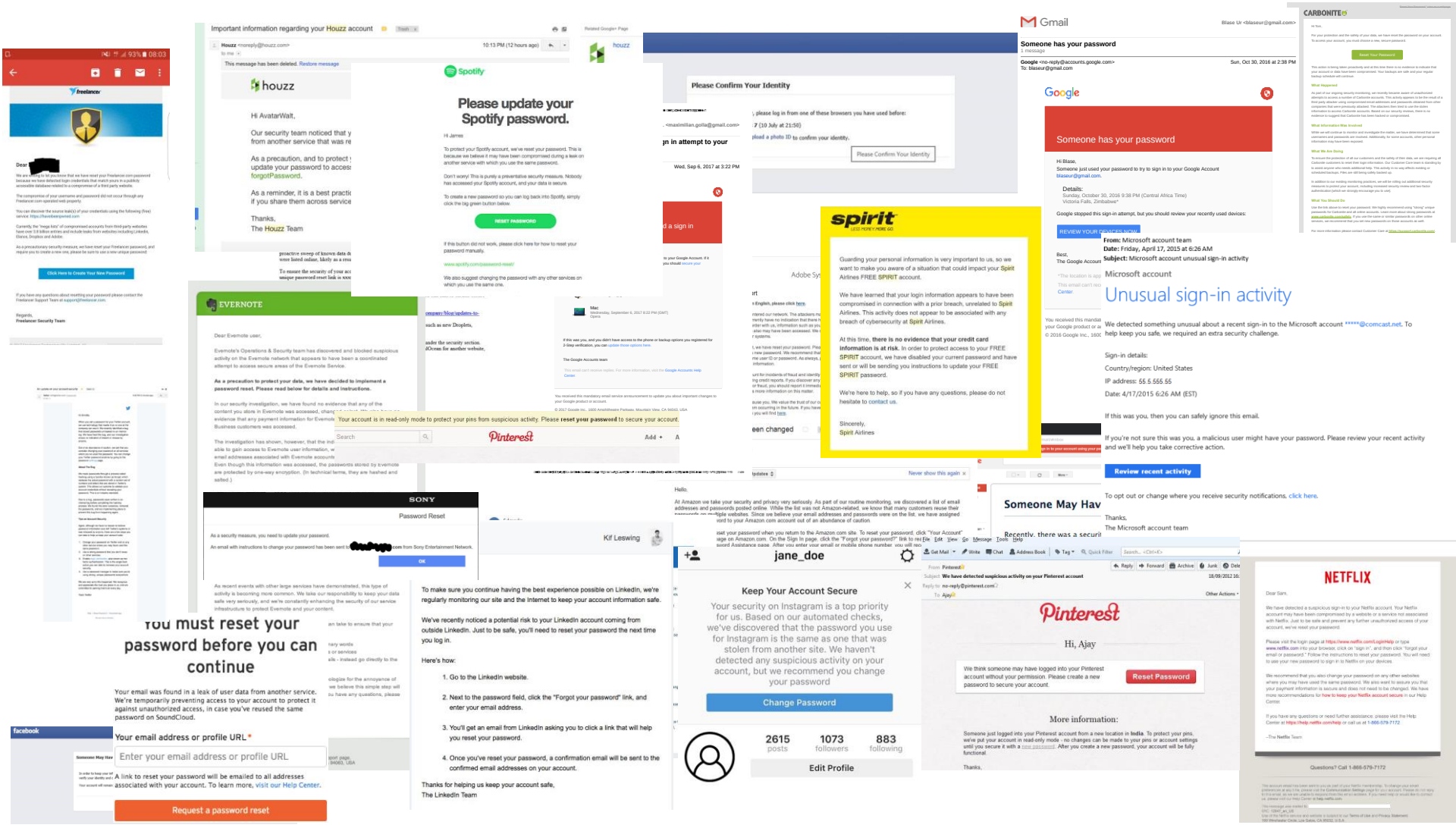
SECURITY

Facebook buys black market passwords to keep your account safe

The company's security chief says account safety is about more than just building secure software.

BY KATIE COLLINS | NOVEMBER 9, 2016 12:56 PM PST

Password-Reuse Notifications



Authentication in Practice: Password Managers

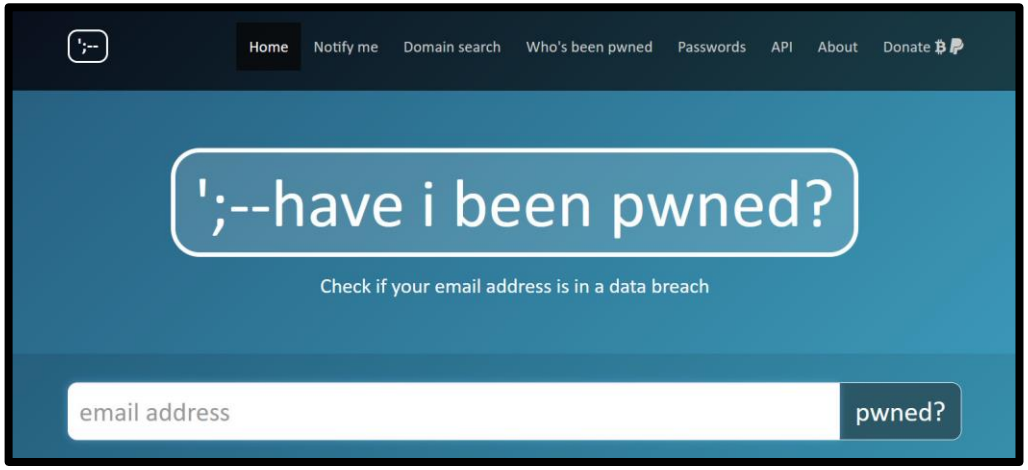
Password Managers

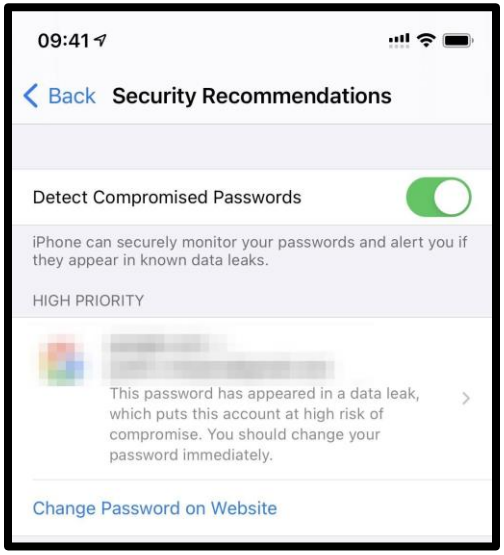
- Trust all passwords to a single master password
  - Also trust software
  - Centralized vs. decentralized architectures

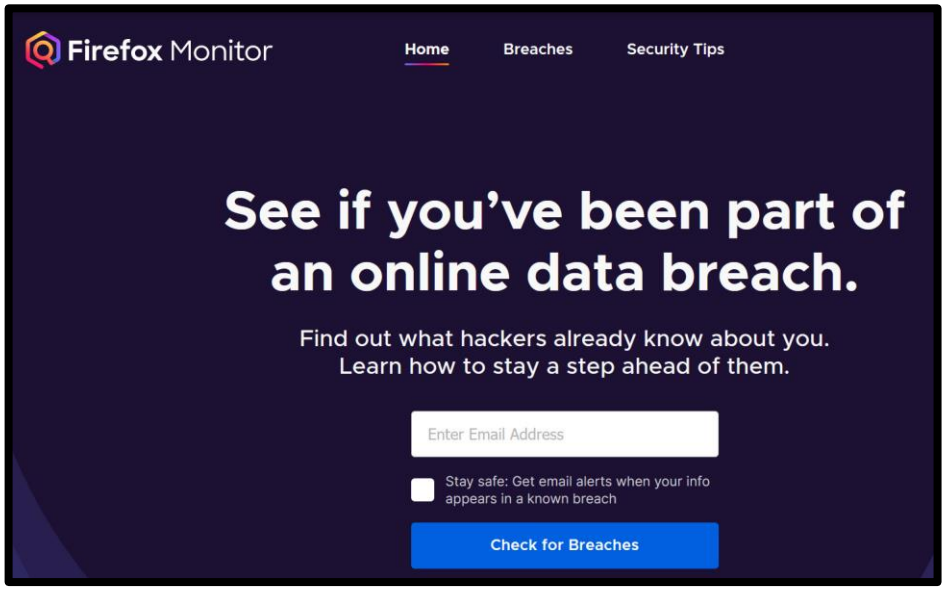
1Password

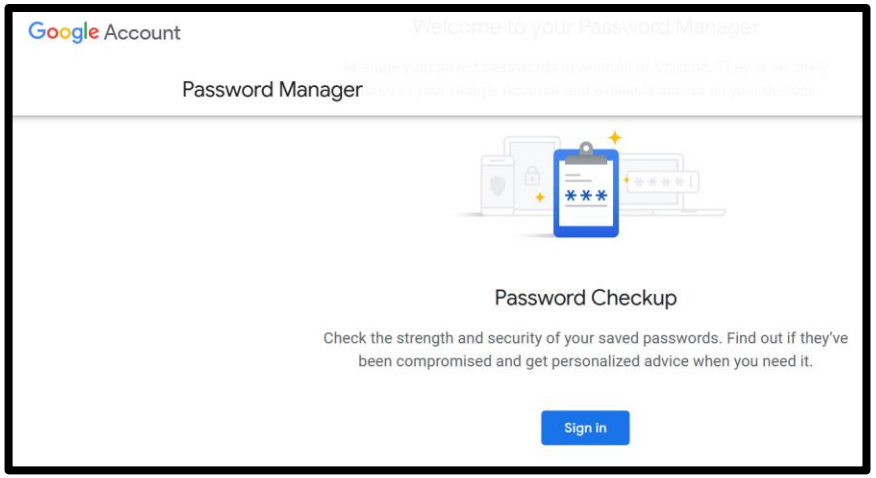
Authentication in Practice: Checking for Compromised Credentials

Checking for Compromised Credentials



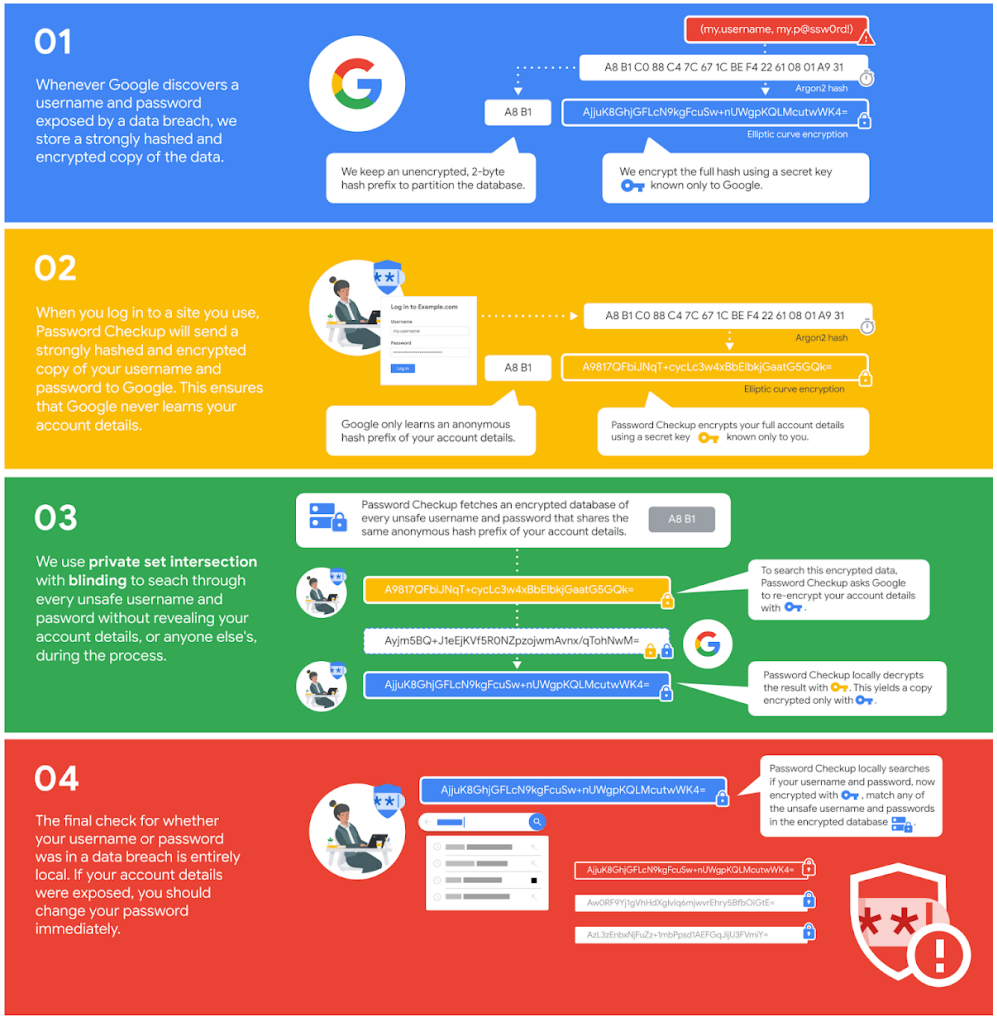






https://www.zdnet.com/article/google-launches-password-checkup-feature-will-add-it-to-chrome-later-this-year/
https://ios.gadgethacks.com/how-to/ios-14-monitors-your-passwords-protect-you-against-data-breaches-heres-works-0341281/

Checking for Compromised Credentials



https://security.googleblog.com/2019/02/protect-your-accounts-from-data.html

What about Biometrics?

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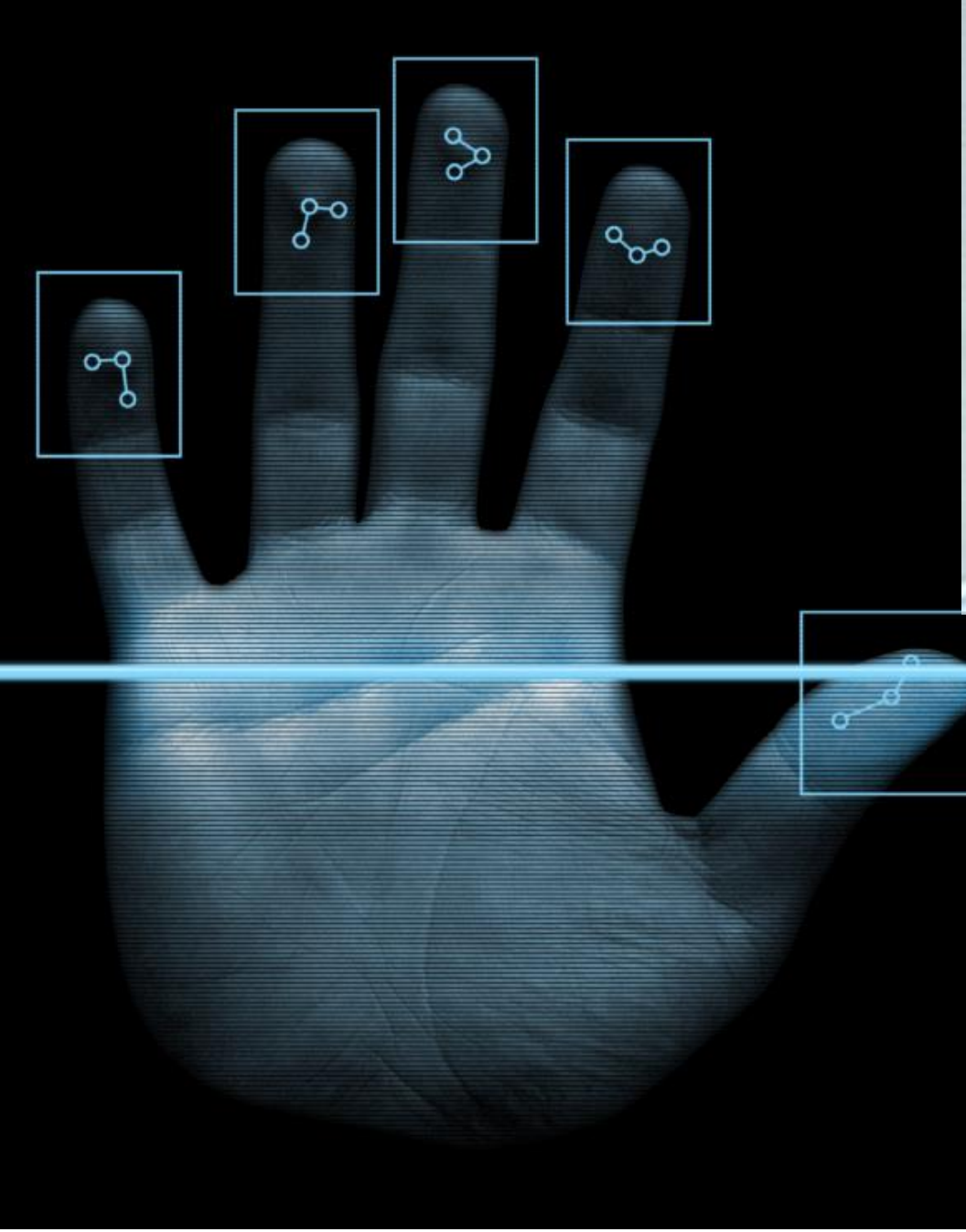

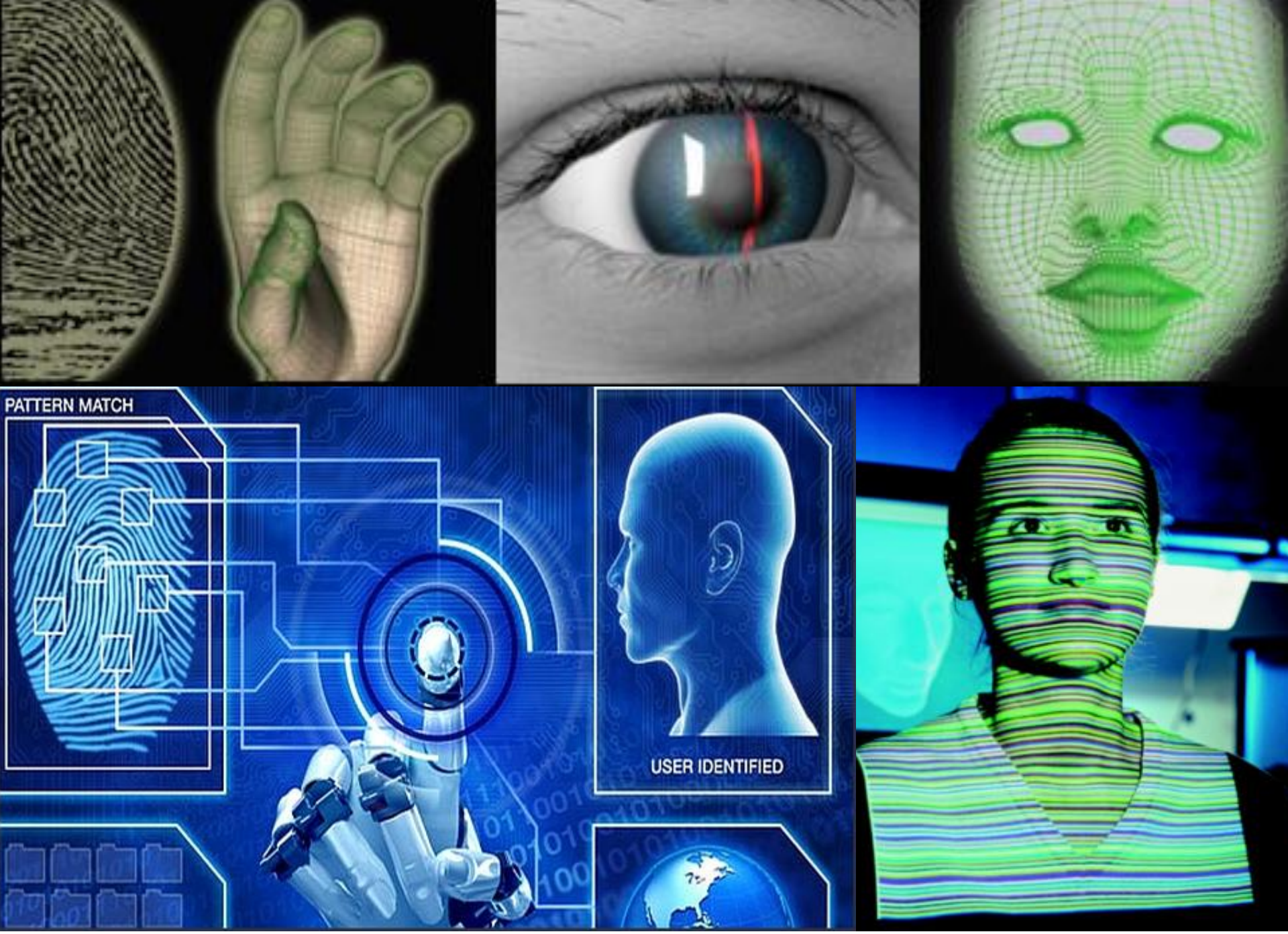

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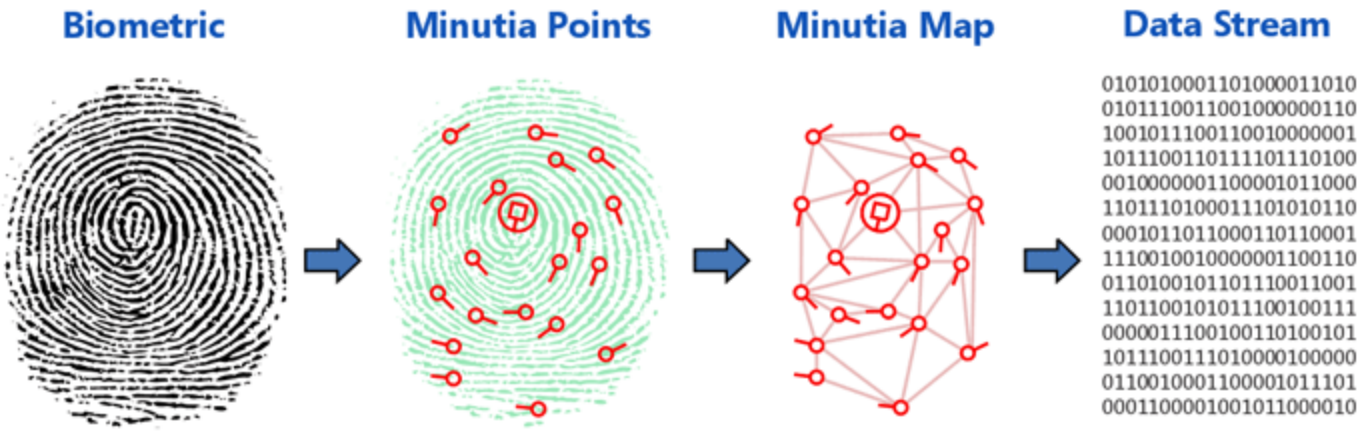
Biometrics

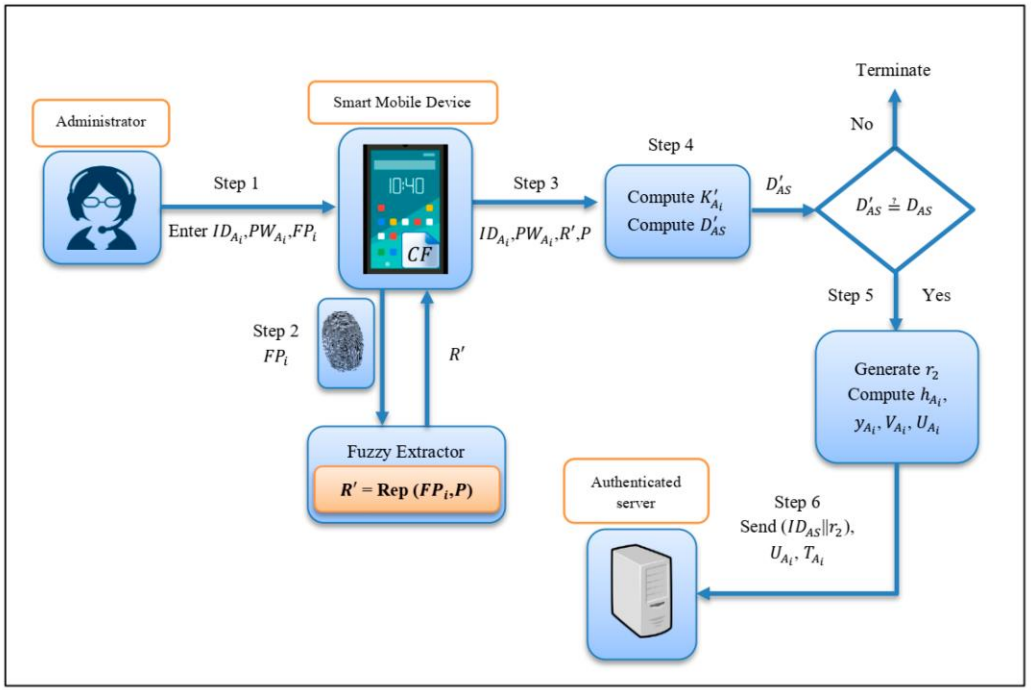
- Fingerprint
- Iris scans or retina scans
- Face recognition
- Finger/hand geometry
- Voice or speech recognition
- The way you type
- (Many others)

Practical Challenges for Biometrics

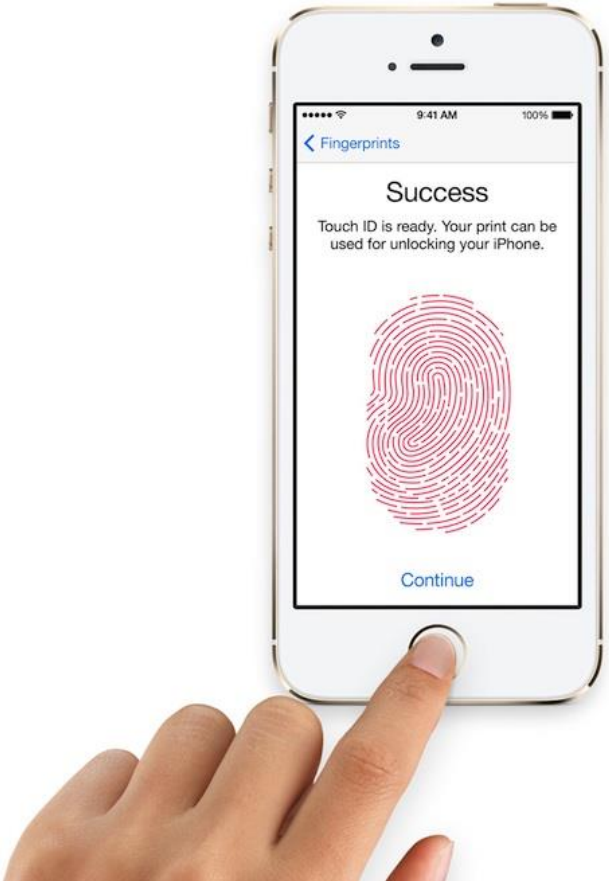
- Immutable (can’t be changed)
- Potentially sensitive data
- High equipment costs
- Sensitive to changes in the environment
- Biometrics can change over time

Storing Biometrics: Templates





iPhone
Touch ID



Android
Face Unlock

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Smartphone Biometrics

- Purpose is to reduce the number of times a user must enter their password
- Falls back to the password
- Face recognition can be tricked by a photo
- Fingerprint recognition can be tricked by a gummy mold
- Users find fingerprint unlock convenient, but do not particularly like face unlock

Conclusions

- Authentication is really hard!
  - Hard for system administrators
  - Hard for users
- Unfortunately, authentication is necessary