

# Negative Numbers

CS143: lecture 15

Byron Zhong, July 20

# What is this?

1110 1001

- 233
- Ascii é
- A pointer pointing to memory location 233 (0xE9)
- Sound wave sample with some magnitude
- A pixel intensity of 91.31%
- ...
- It's one of the 256 choices, whatever that is.

# Negative Numbers

- If we have 32 bits, we have  $2^{32}$  choices
  - unsigned:  $2^{32}$  choices of non-negative numbers -- 0 through  $2^{32} - 1$
  - int:  $2^{32}$  choices of both negative and positive numbers *about*  $-2^{31}$  through  $2^{31}$

# Negative Numbers

## Attempt 1

- The first bit is the sign, 0 -> positive, 1 -> negative

5 = 0000 0000 0000 0000 0000 0000 0000 0011

-5 = 1000 0000 0000 0000 0000 0000 0000 0011

- Range?
  - $[-(2^{31} - 1), 2^{31} - 1]$

# Negative Numbers

## Attempt 1

- The first bit is the sign, 0 -> positive, 1 -> negative

5 = 0000 0000 0000 0000 0000 0000 0000 0011

-5 = 1000 0000 0000 0000 0000 0000 0000 0011

- What's wrong with this?

0000 0000 0000 0000 0000 0000 0000 0000

1000 0000 0000 0000 0000 0000 0000 0000

# Negative Numbers

## Attempt 1

- The first bit is the sign, 0 -> positive, 1 -> negative

5 = 0000 0000 0000 0000 0000 0000 0000 0011

-5 = 1000 0000 0000 0000 0000 0000 0000 0011

- What's wrong with this?

0000 0000 0000 0000 0000 0000 0000 0000

1000 0000 0000 0000 0000 0000 0000 0000

- We have +0 and -0 !?
  - Waste a combination, equality check is complicated +0 == -0

# Negative Numbers

## Attempt 1

- The first bit is the sign, 0 -> positive, 1 -> negative

5 = 0000 0000 0000 0000 0000 0000 0000 0011

-5 = 1000 0000 0000 0000 0000 0000 0000 0011

- What's wrong with this?

0000 0000 0000 0000 0000 0000 0000 0011

+ 1000 0000 0000 0000 0000 0000 0000 0011

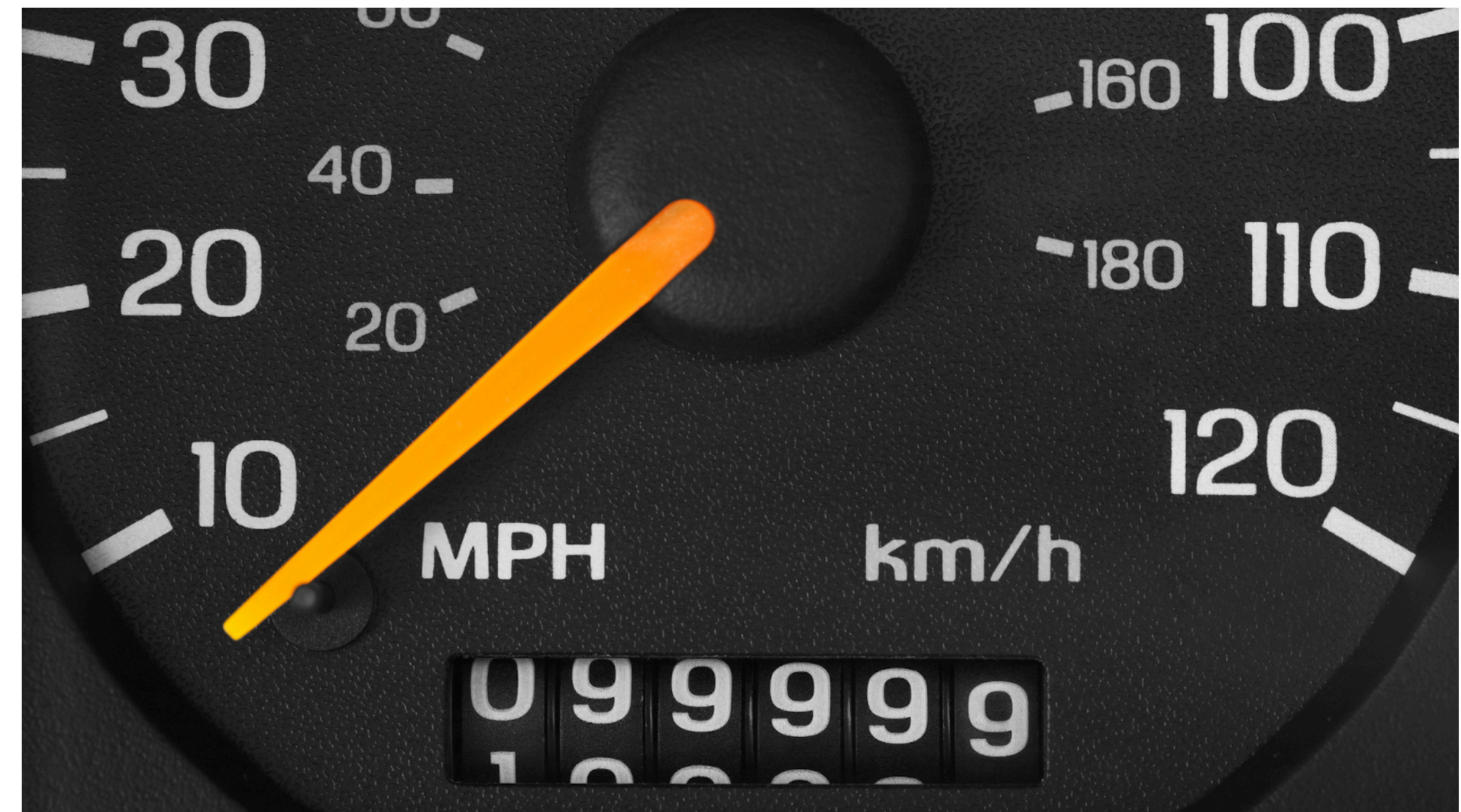
- 5 + (-5) is not 0 :(((

- Need special circuit for adding signed numbers and unsigned numbers



# Negative Numbers

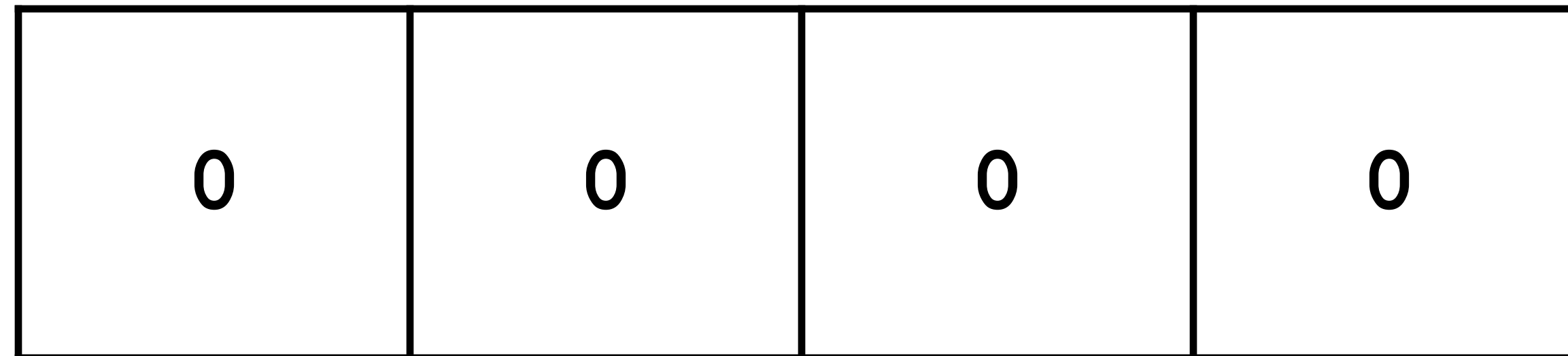
**Better Idea: Two's Complement**





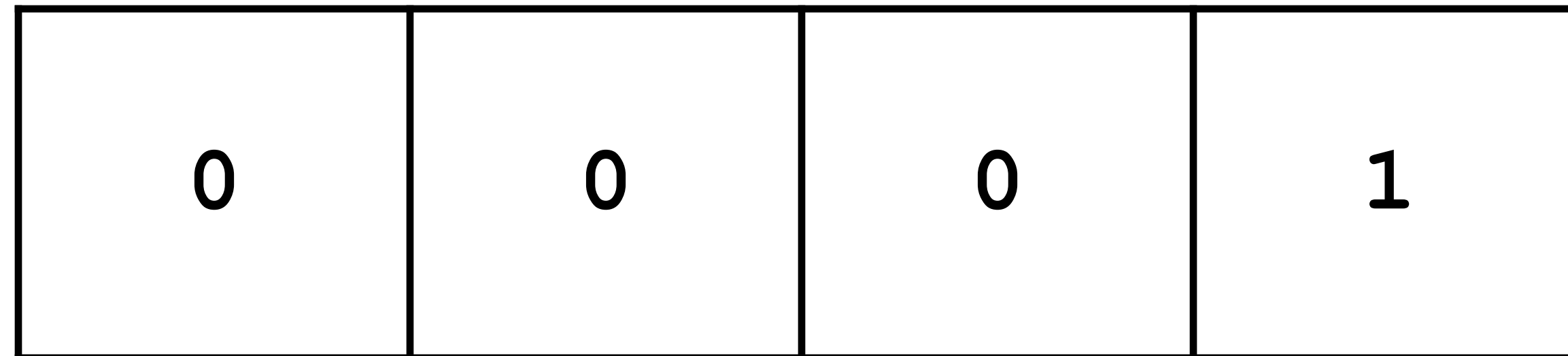
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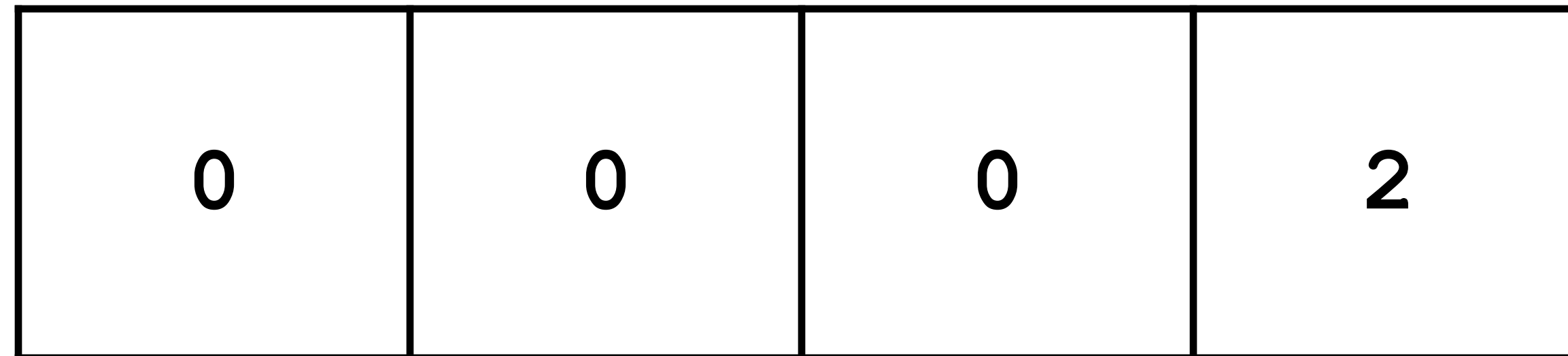
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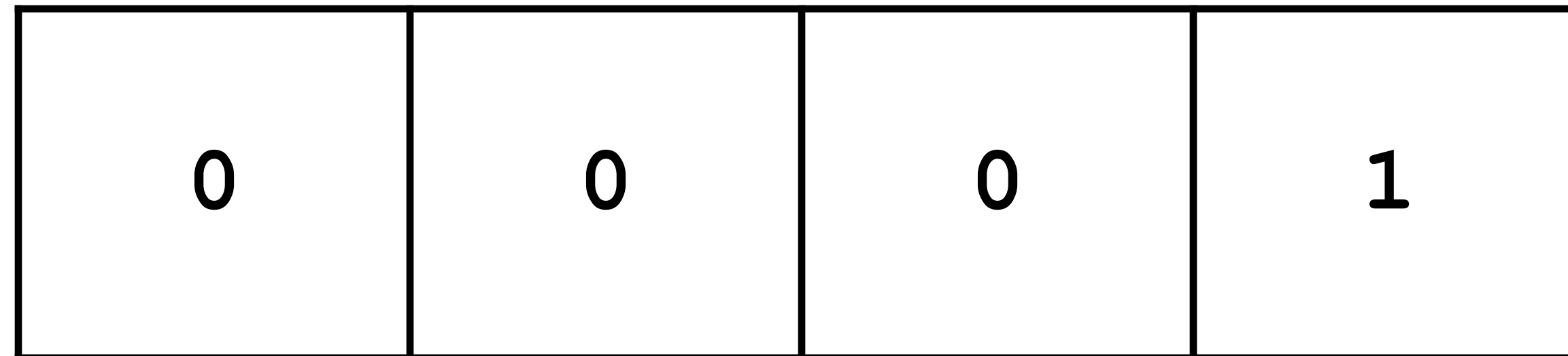
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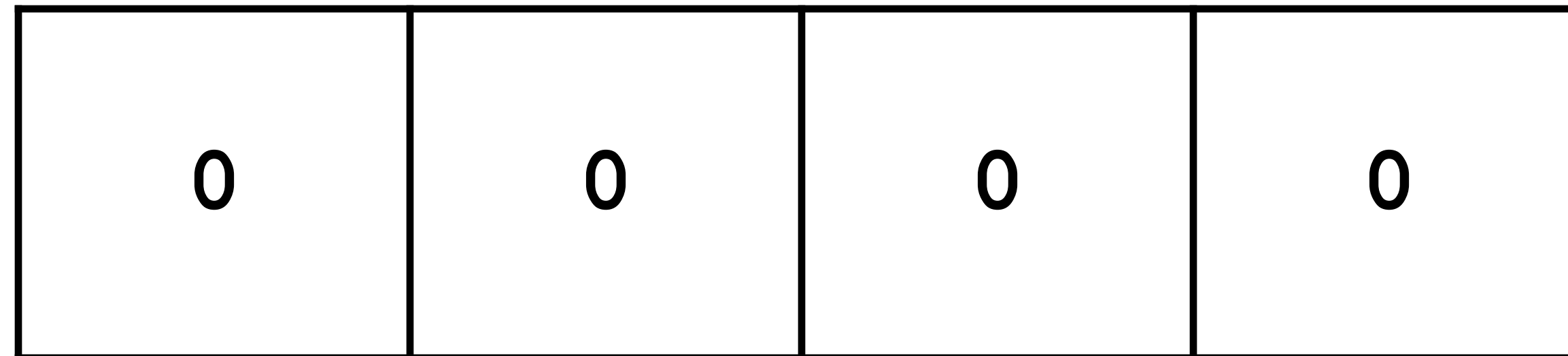
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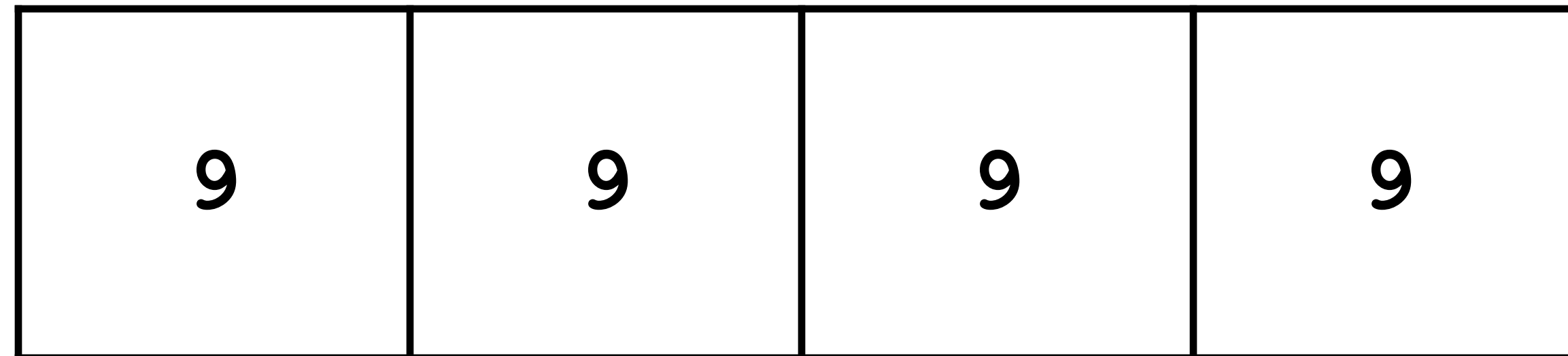
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# Negative Numbers

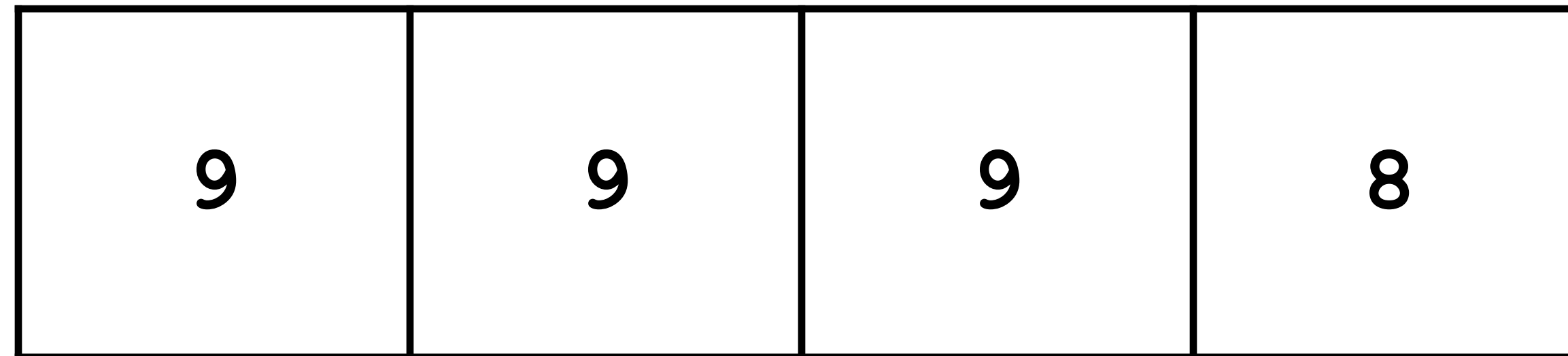
**Better Idea: Two's Complement**





# Negative Numbers

**Better Idea: Two's Complement**



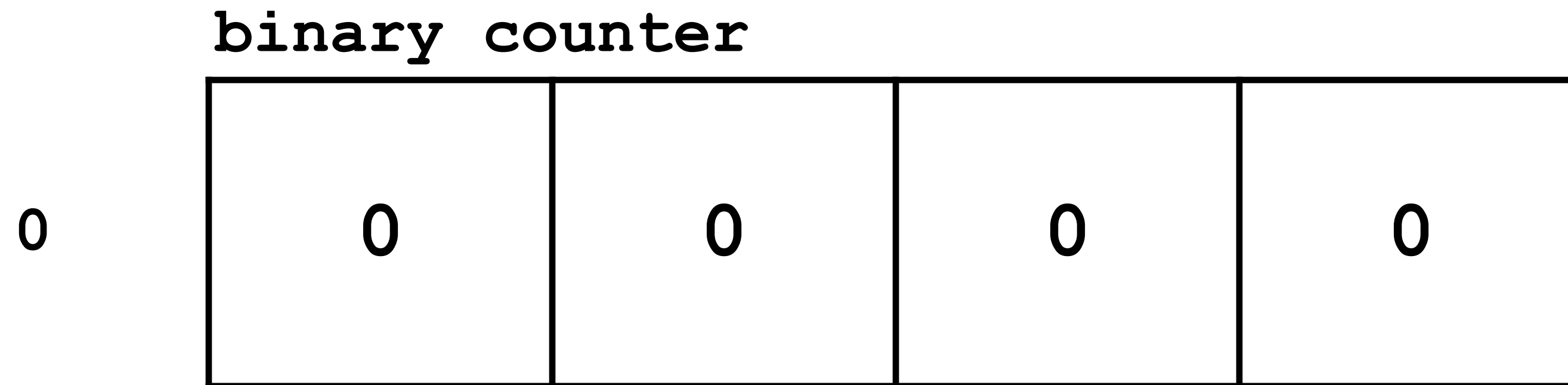
# Negative Numbers

**Better Idea: Two's Complement**

9	9	9	7
---	---	---	---

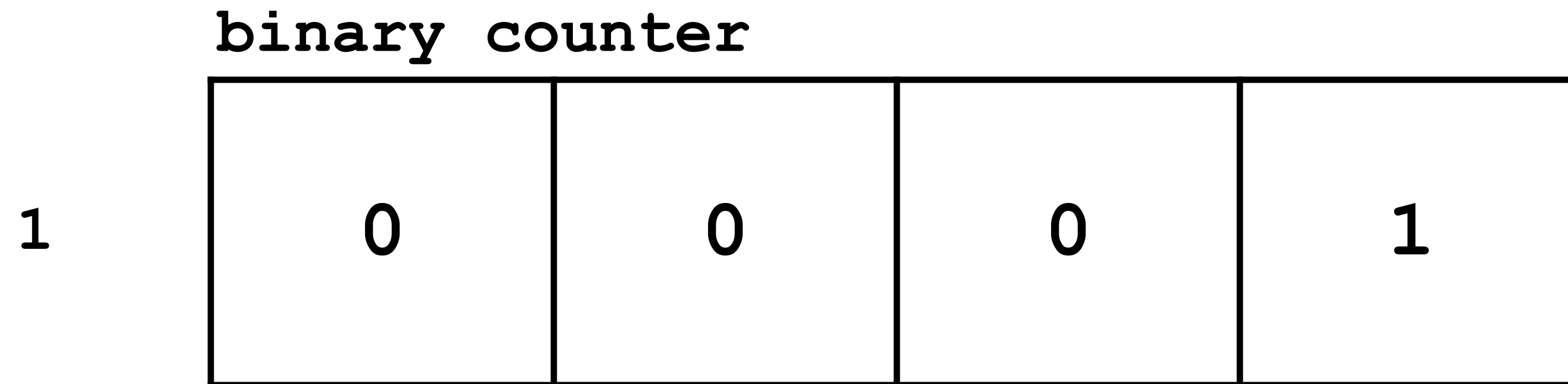
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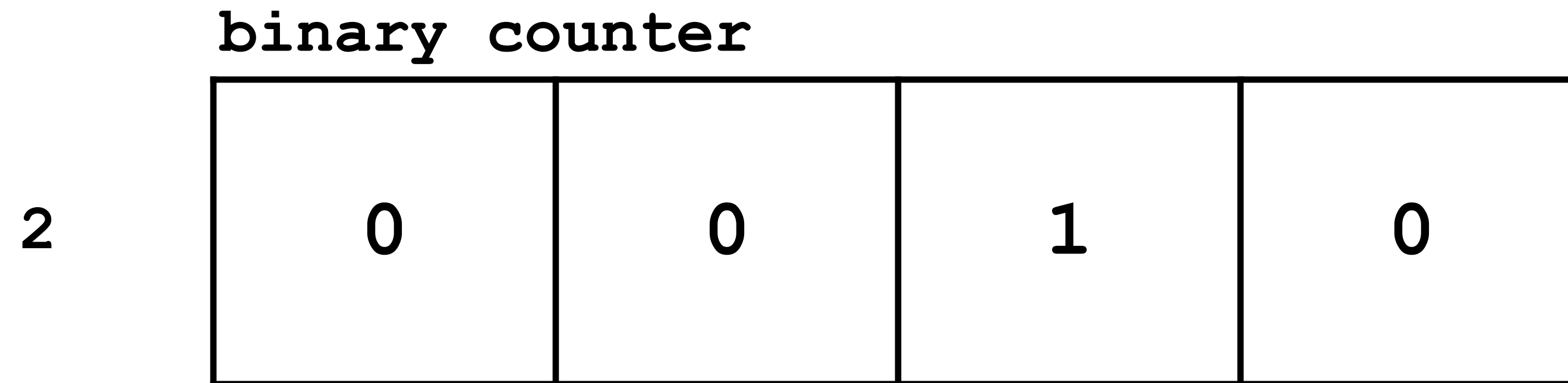
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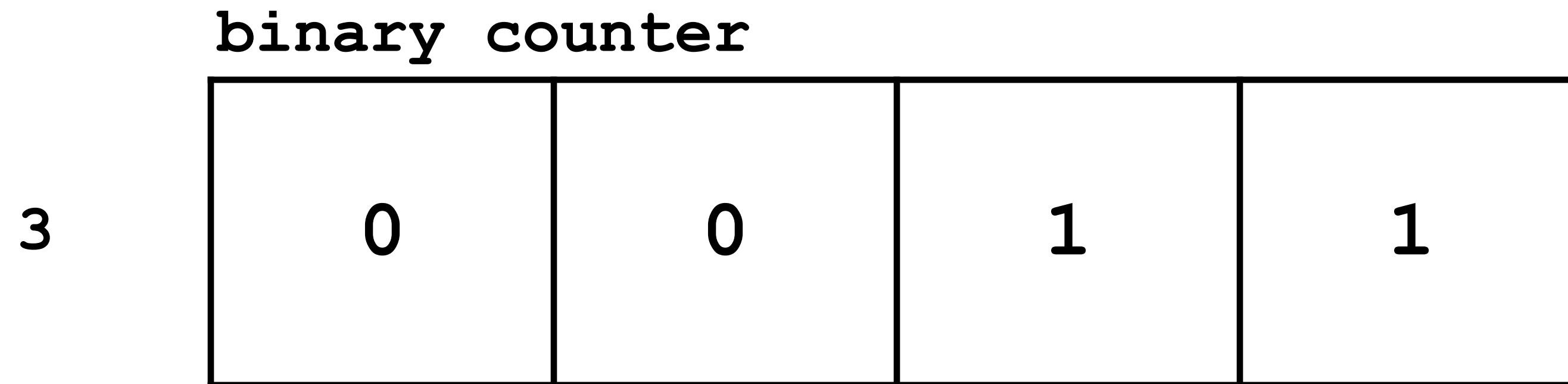
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**Better Idea: Two's Complement**



# Negative Numbers

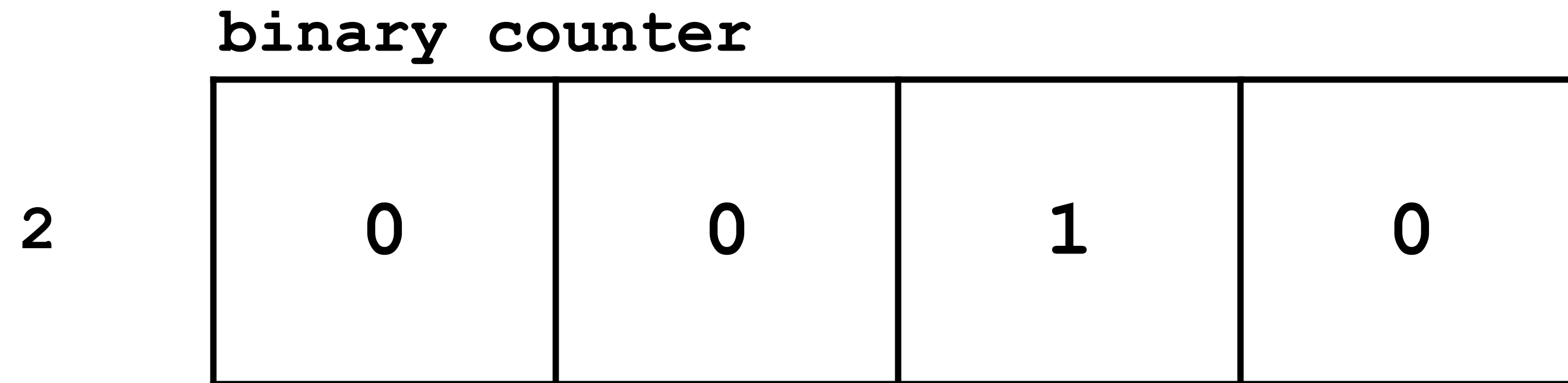
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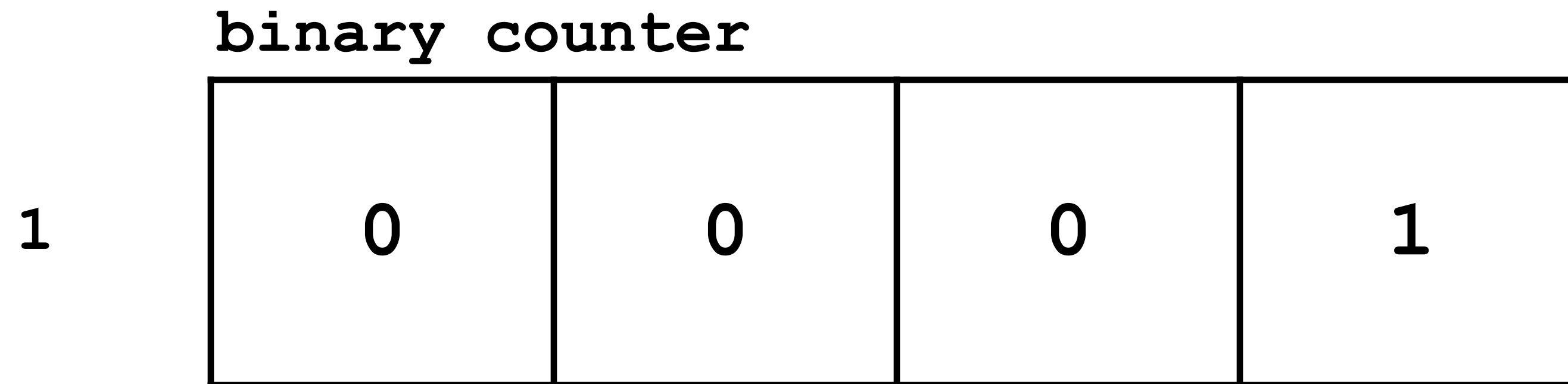
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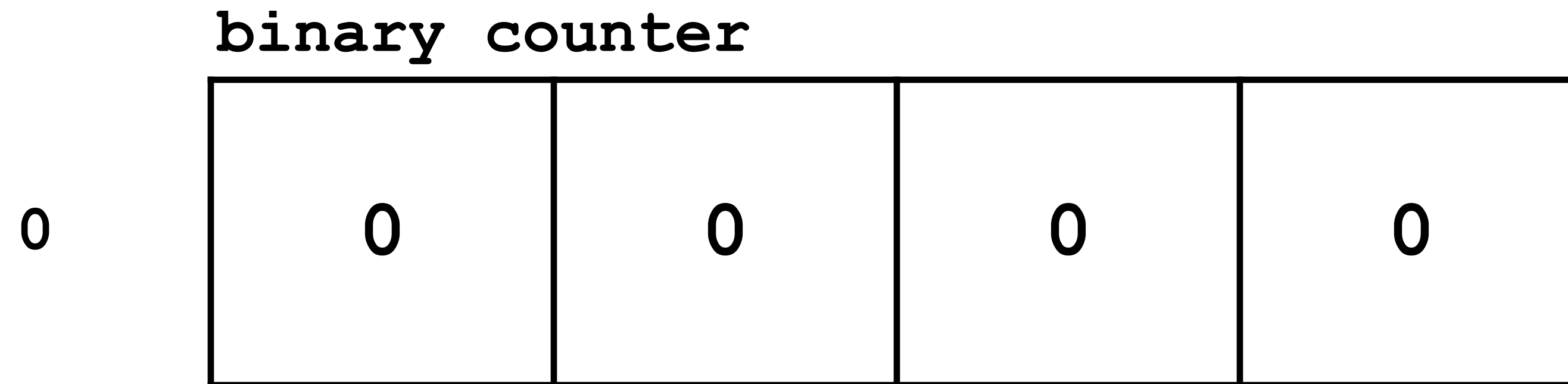
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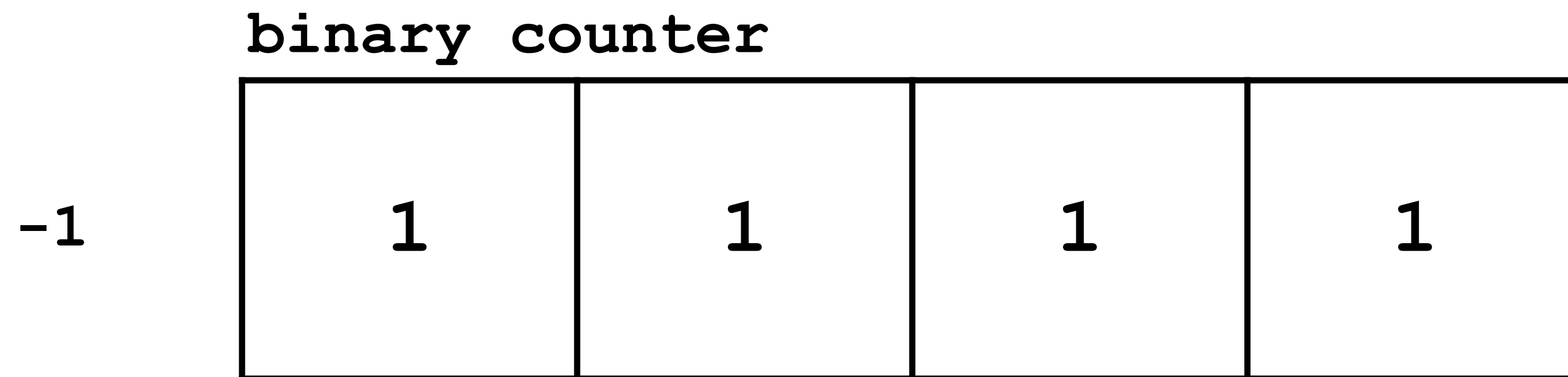
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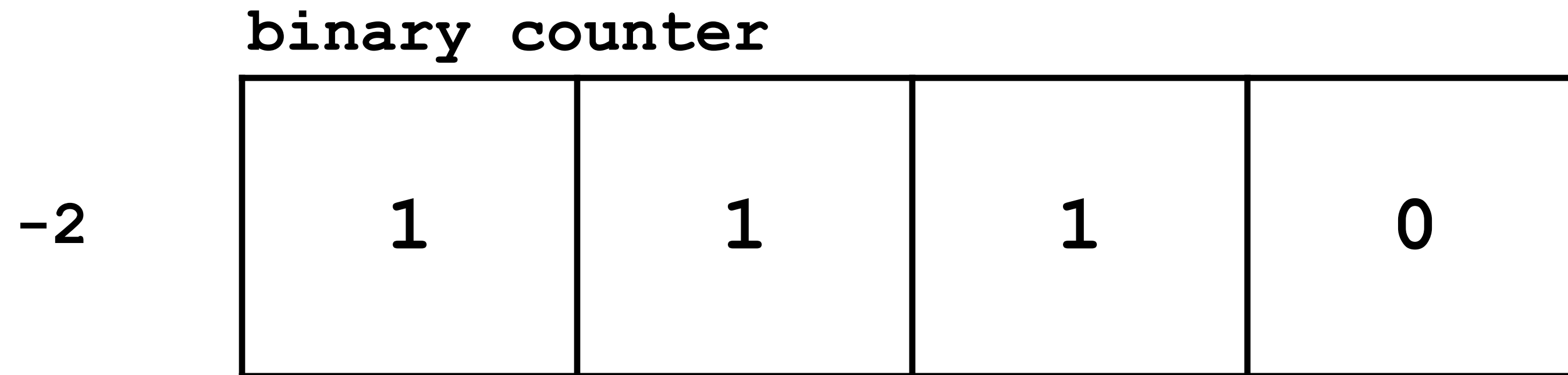
# Negative Numbers

Better Idea: Two's Complement



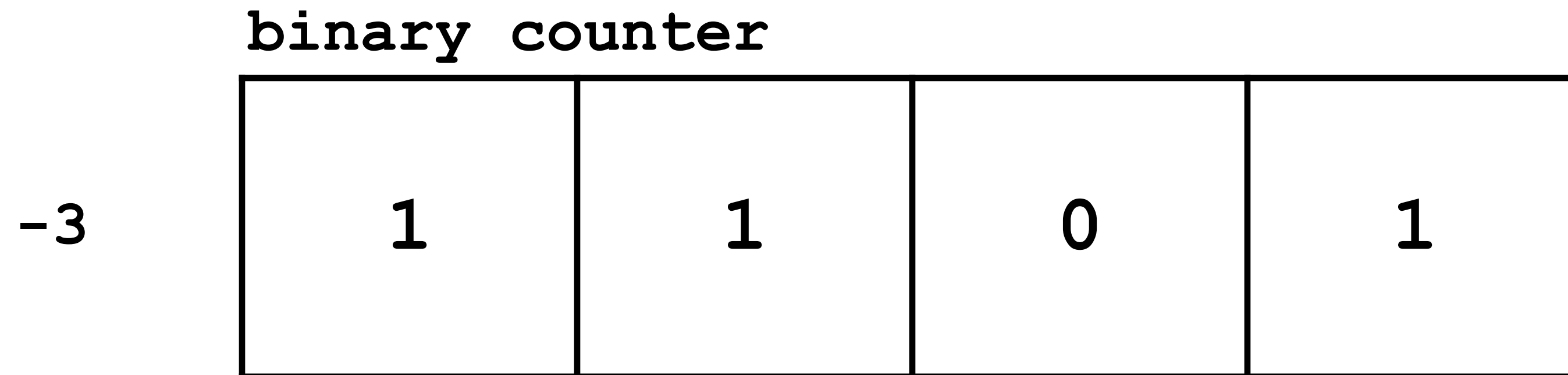
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# Negative Numbers

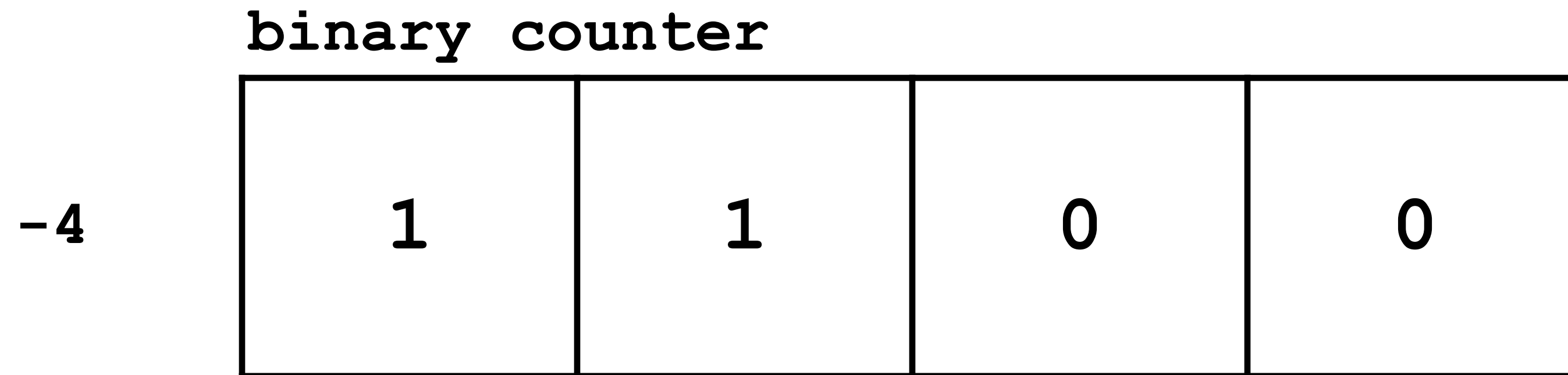
Better Idea: Two's Complement





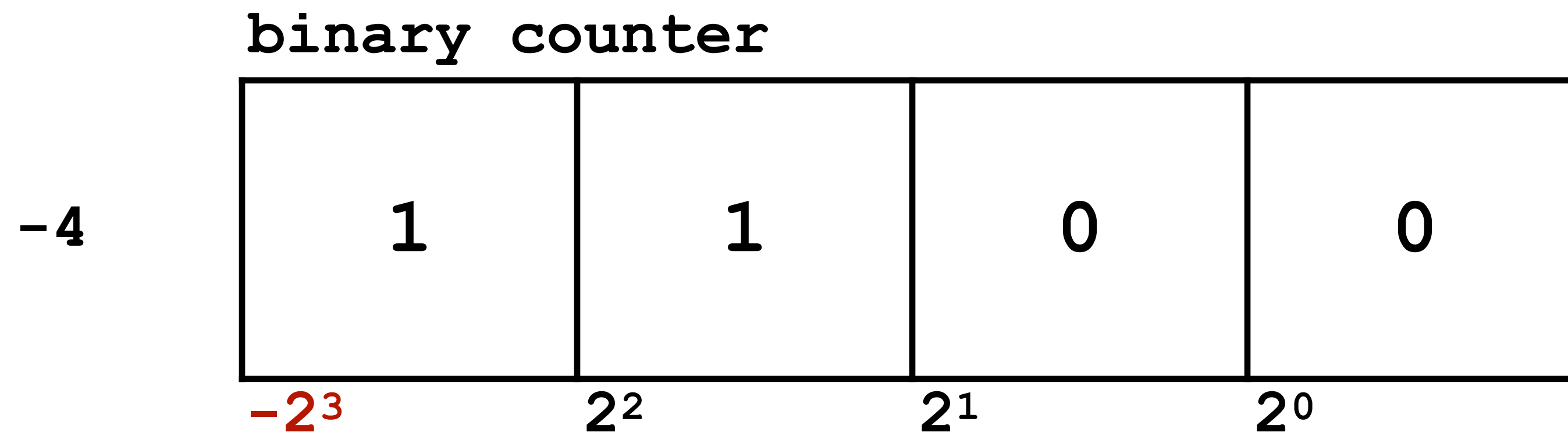
# Negative Numbers

Better Idea: Two's Complement



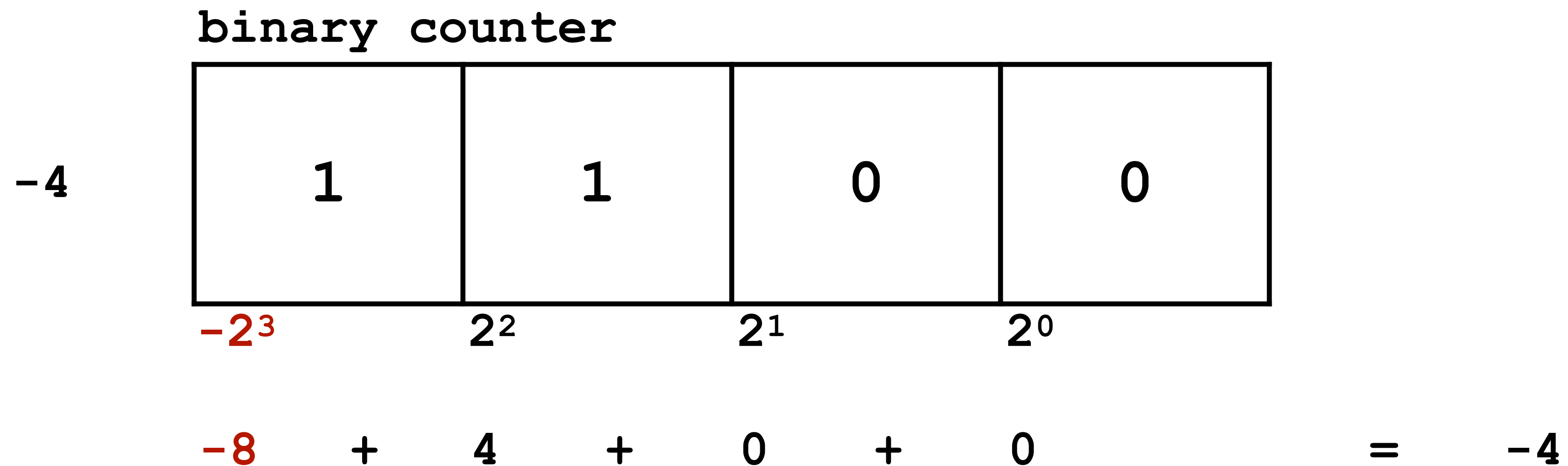
# Negative Numbers

## Two's complement: two ways



# Negative Numbers

## Two's complement: two ways



- The most significant bit is negative.

# Negative Numbers

## Two's complement: two ways

n	2^n
0	1
1	2
2	4
3	8
4	16
5	32
6	64
7	128

1110 1001

- Unsigned: 233
- What does this represent if it is a signed integer?
- $-128 + 64 + 32 + 8 + 1 = -23$

# Negative Numbers

## Two's complement: two ways

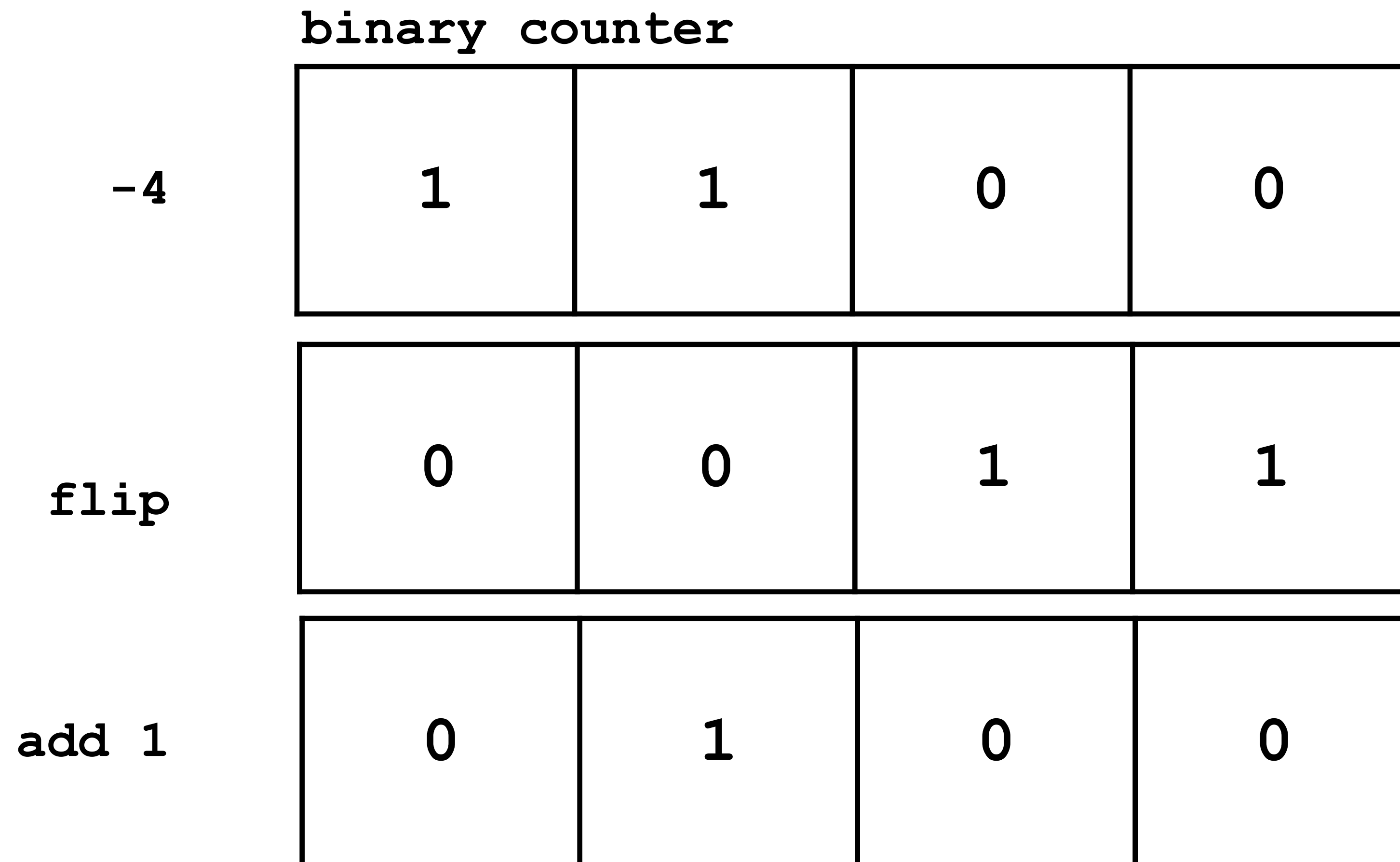
n	2^n
0	1
1	2
2	4
3	8
4	16
5	32
6	64
7	128

1110 1001

- How do hardwares negate the sign?
  - Flip all the bits and add one
- Flip all the bits: 0001 0110
- add one: 0001 0111
- : 16 + 4 + 2 + 1 = 23

# Negative Numbers

## Two's complement: two ways





# Negative Numbers

Two's complement: range

INT\_MAX

1	0000	0001
0	0000	0000
-1	1111	1111
-2	1111	1110

...

...

# Negative Numbers

Two's complement: range

INT\_MAX ↑

0111 1111

...

1 0000 0001

0 0000 0000

-1 1111 1111

-2 1111 1110

...

# Negative Numbers

Two's complement: range

INT\_MAX ↑

0111 1111

...

1 0000 0001

0 0000 0000

-1 1111 1111

-2 1111 1110

...

INT\_MIN

# Negative Numbers

Two's complement: range

INT\_MAX ↑ 0111 1111

...

1 0000 0001

0 0000 0000

-1 1111 1111

-2 1111 1110

...

INT\_MIN | 1000 0000

# Negative Numbers

Two's complement: range

INT\_MAX | 0111 1111 <---  $2^7 - 1$

...

1 | 0000 0001

0 | 0000 0000

-1 | 1111 1111

-2 | 1111 1110

...

INT\_MIN | 1000 0000 <---  $-2^7$

# Negative Numbers

Two's complement: range

INT\_MAX | 0111 1111 <---  $2^8 - 1$

...

1 | 0000 0001

0 | 0000 0000

-1 | 1111 1111

-2 | 1111 1110

...

INT\_MIN | 1000 0000 <---  $-2^8$

# Signed Arithmetics

INT_MAX	↑	0111	1111
		...	
1		0000	0001
0		0000	0000
-1		1111	1111
-2		1111	1110
		...	
INT_MIN		1000	0000

$$\begin{array}{r} 1111\ 1110 \\ + 0000\ 0011 \\ \hline \end{array}$$

# Signed Arithmetics

INT_MAX	0111	1111
	...	
1	0000	0001
0	0000	0000
-1	1111	1111
-2	1111	1110
	...	
INT_MIN	1000	0000

$$\begin{array}{r} \phantom{+} 1111 \ 1110 \\ + 0000 \ 0011 \\ \hline \end{array} \quad \begin{array}{r} -2 \\ 3 \end{array}$$



# Signed Arithmetics

INT_MAX	0111	1111
	...	
1	0000	0001
0	0000	0000
-1	1111	1111
-2	1111	1110
	...	
INT_MIN	1000	0000

$$\begin{array}{r} \phantom{+} 1111 \ 1110 \\ + 0000 \ 0011 \\ \hline \phantom{+} \phantom{1111} \ 1 \end{array}$$

-2  
3

# Signed Arithmetics

INT_MAX	0111	1111
	...	
1	0000	0001
0	0000	0000
-1	1111	1111
-2	1111	1110
	...	
INT_MIN	1000	0000

$$\begin{array}{r} \phantom{+} 1111 \ 1110 \\ + 0000 \ 0011 \\ \hline \phantom{+} \phantom{0000} \ 01 \end{array}$$

-2  
3

# Signed Arithmetics

INT_MAX	0111	1111
	...	
1	0000	0001
0	0000	0000
-1	1111	1111
-2	1111	1110
	...	
INT_MIN	1000	0000

$$\begin{array}{r} \phantom{+} 1111 \ 1110 \\ + 0000 \ 0011 \\ \hline \phantom{+} \phantom{0000} \ 001 \end{array}$$

-2  
3

# Signed Arithmetics

INT_MAX	0111	1111
	...	
1	0000	0001
0	0000	0000
-1	1111	1111
-2	1111	1110
	...	
INT_MIN	1000	0000

$$\begin{array}{r} \phantom{+} 1111 \ 1110 \\ + 0000 \ 0011 \\ \hline \phantom{+} \phantom{0000} 0001 \end{array} \quad \begin{array}{l} -2 \\ 3 \end{array}$$

# Signed Arithmetics

INT_MAX	0111	1111
	...	
1	0000	0001
0	0000	0000
-1	1111	1111
-2	1111	1110
	...	
INT_MIN	1000	0000

$$\begin{array}{r} 1111\ 1110 \\ + 0000\ 0011 \\ \hline 0\ 0001 \end{array}$$

-2  
3

# Signed Arithmetics

INT_MAX	0111	1111
	...	
1	0000	0001
0	0000	0000
-1	1111	1111
-2	1111	1110
	...	
INT_MIN	1000	0000

$$\begin{array}{r} \phantom{+} 1111 \ 1110 \\ + 0000 \ 0011 \\ \hline \phantom{+} 00 \ 0001 \end{array} \quad \begin{array}{l} -2 \\ 3 \end{array}$$

# Signed Arithmetics

INT_MAX	0111	1111
	...	
1	0000	0001
0	0000	0000
-1	1111	1111
-2	1111	1110
	...	
INT_MIN	1000	0000

$$\begin{array}{r} 1111\ 1110 \\ +\ 0000\ 0011 \\ \hline 000\ 0001 \end{array}$$

-2  
3

# Signed Arithmetics

INT_MAX	↑	0111	1111
		...	
1		0000	0001
0		0000	0000
-1		1111	1111
-2		1111	1110
		...	
INT_MIN		1000	0000

$$\begin{array}{r} 1111\ 1110 \quad -2 \\ + 0000\ 0011 \quad 3 \\ \hline 0000\ 0001 \end{array}$$

- The hardware does not need to know/care whether the arithmetics is signed or unsigned
- $\text{INT\_MAX} + 1 == \text{INT\_MIN}$



# Signed Arithmetics

Unsigned >>

```
      1111 1111 1111 1111 1111 1111 1111 1100      -4  
>>                                     1  
-----
```

# Signed Arithmetics

## Unsigned >>

```
      1111 1111 1111 1111 1111 1111 1111 1100          -4
    >>
-----
    0111 1111 1111 1111 1111 1111 1111 1110      2147483646
```

- If we just add 0 in right shift, shifting a negative number results in a positive number.
- Ideally, we would like to keep the sign / shifting is dividing by 2.
- We repeat the most significant bit when doing a *signed* right shift.

# Signed Arithmetics

Signed >>

```
          1111 1111 1111 1111 1111 1111 1111 1100          -4
>>                                     1
-----
      1111 1111 1111 1111 1111 1111 1111 1110
```

# Signed Arithmetics

## Signed >>

	1111 1111 1111 1111 1111 1111 1111 1100	-4
>>		1
<hr/>		
	1111 1111 1111 1111 1111 1111 1111 1110	-2
	0000 0000 0000 0000 0000 0000 0000 0100	4
>>		1
<hr/>		

# Signed Arithmetics

## Signed >>

	1111 1111 1111 1111 1111 1111 1111 1100	-4
>>		1
<hr/>		
	1111 1111 1111 1111 1111 1111 1111 1110	-2
	0000 0000 0000 0000 0000 0000 0000 0100	4
>>		1
<hr/>		
	0000 0000 0000 0000 0000 0000 0000 0010	2

- If the first bit is 0, right shifting fills with 0.
- If the first bit is 1, right shifting fills with 1.

# Signed Arithmetics

## Signed >>

- In C, *signed* right shift and *unsigned* right shift looks exactly the same
- C will use *signed* shift if the variable being shifted is a *signed* integer

```
int x = ~0;  
x = x >> 4;  
printf("0x%x\n", x);
```

0xffffffff

```
unsigned int x = ~0;  
x = x >> 4;  
printf("0x%x\n", x);
```

0x0fffffff

# Signed Arithmetics

- Signed extension works the same way:

```
short x = ~0;  
int y = x;  
printf("0x%x\n", y);
```

0xffffffff

```
unsigned short x = ~0;  
unsigned int y = x;  
printf("0x%x\n", y);
```

0x0000ffff

# Machine Arithmetics

## Summary

- For a 32-bit integer:
  - Range: unsigned  $[0, 2^{32} - 1]$ ; signed  $[-2^{31}, 2^{31} - 1]$ .
  - Meaning of the highest bit: unsigned  $2^{31}$ , signed  $-2^{31}$ .
- Negate a number: flips all bits and add 1 ( $-n == \sim n + 1$ ).
- Most arithmetic operators work regardless of the signedness of the integer except right shift:
  - Signed right shift (arithmetic right shift): fills with the highest bit.
  - Unsigned right shift (logical right shift): fills with 0.
  - Left shift is the same.

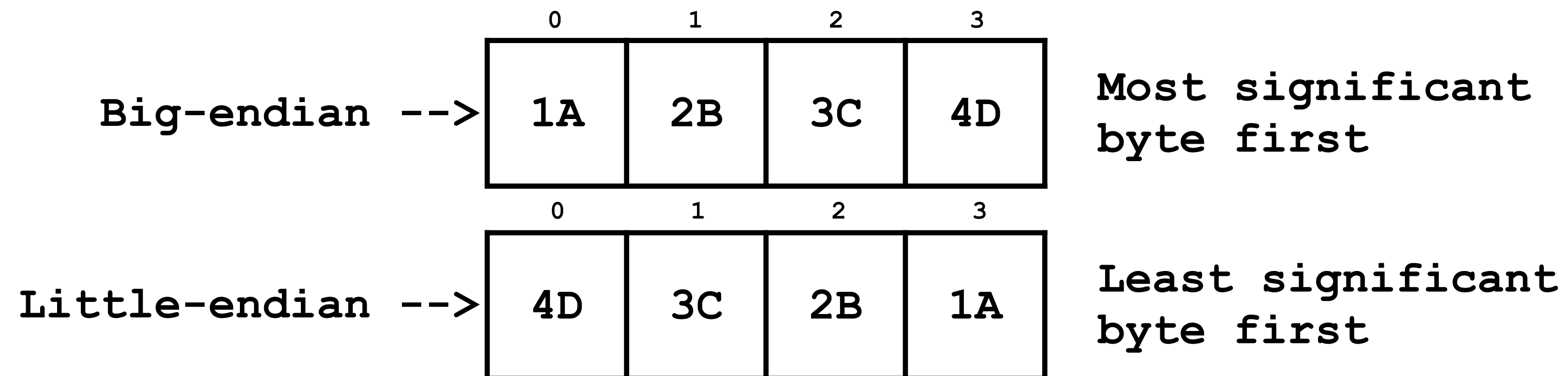


# Enum

## Demo

# Endian

- We think of an integer as one atomic value:
  - `int x = 0x1A2B3C4D;`
- But if an integer has 4 bytes and each byte is addressable, which of the 4 bytes is stored first?



# Endian

- Is my machine little-endian or big-endian?
- Let's find out!

# Endian

- Our machine is little-endian?????
- We usually write numbers in big-endian: 345 is three hundred and forty-five
- But there are some advantages for little-endian:
  - comparing two numbers of different length (long and int e.g.)
    - 4E3C2B1A
    - 4E3C2B1A00000000
  - addition, subtraction circuits work from low to high
  - etc.

# Endian

## Does it matter?

- Mostly we don't care. Unless you do memory trickery, variables work as you would expect
- However, when we serialize data into byte sequences, you need to pay extra attention:
  - Writing a number to a file
  - Sending a number over a network
- You and the reader must agree on byte order
  - For this purpose, *network byte order* is defined for TCP/IP

# Octal Number



# Octal Number

- Base-8 numbers, 0-7
- Corresponds to 3 bits
- In C, prefix by 0, e.g. `0123`, is 83
- Not very common

# Bits

- A bit answers a yes/no question
  - Bits have no inherent meanings; to read encoding, need to know the intended interpretation
  - Any device that can be one of two states can encode a bit.
- Binary and hexadecimal numbers
  - $xyz_m = x \cdot m^2 + y \cdot m^1 + z \cdot m^0$
  - One hex digit corresponds two 4 binary digits (bits) -- easy conversion between the two
- Positive and negative numbers
- Arithmetics
  - +, -, <<, &, |, ^, ~, two >>, negate
  - Bit-packing
- Endianness